

The y team

(joint with Armin Straub) 01/12/2023

$$A^{(s)}(n) = \sum_{k=0}^n \binom{n}{k}^s \quad \text{generalized Franel numbers}$$

$$A^{(1)}(n) = 2^n \quad A^{(1)}(n+1) - 2A^{(1)}(n) = 0$$

$$A^{(2)}(n) = \binom{2n}{n} \quad (n+1)A^{(2)}(n+1) - 2(2n+1)A^{(2)}(n) = 0$$

Franel 1894 $(n+1)^2 A^{(3)}(n+1) - (7n^2 + 7n + 2)A^{(3)}(n) - 8n^2 A^{(3)}(n-1) = 0$

Franel 1895 $(n+1)^3 A^{(4)}(n+1) - 2(2n+1)(3n^2 + 3n + 1)A^{(4)}(n)$

$$- 4n(4n-1)(4n+1)A^{(4)}(n-1) = 0$$

Predicted a general shape of the recursion,
in particular that the order is $\lfloor \frac{s+1}{2} \rfloor$

1987 Perlsstadt: $s \geq 5, 6$

1989 McIntosh: $s \geq 5, 6, \dots, 10$

Nowadays (3, 1990) such things are done using the algorithm of creative telescoping invented by Zeilberger

What it does?

telescoper
for $a(n, k)$

It takes a hypergeometric term $a(n, k)$, e.g. $\binom{n}{k}^s$,
and returns the difference operator $P(n, N) \in \mathbb{Z}[n][N]$,
where $N: n \mapsto n+1$, and another hypergeometric
term $b(n, k)$ such that

$$P(n, N)a(n, k) = b(n, k+1) - b(n, k)$$

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Sum both sides over $k \in \mathbb{Z}$ (assuming, for example, that the support in k , for any fixed n is finite),
 of $a(n, k)$ & $b(n, k)$

$$P(n, N) \sum_{k \in \mathbb{Z}} a(n, k) = 0$$

A caveat: the operator $P(n, N)$ may be not of minimal order (in N), reducible.

For example, when $\sum_{k \in \mathbb{Z}} a(n, k)$ admits a different representation, $\sum_{k \in \mathbb{Z}} \tilde{a}(n, k)$ say.

Then $\tilde{P}(n, N)$ and $\tilde{a}(n, k)$ can be different from $P(n, N)$ while $\tilde{a}(n, k)$ is certainly different.

For example, run on $A^{(3)}(n) = \sum_{k=0}^n \binom{n}{k}^2 \binom{2k}{n}$
 to get $\tilde{P} = P$ but different $\tilde{a}(n, k)$.

Conclusion: $P(n, N)$ and its minimality may depend on the choice of $a(n, k)$.

Our result with Armin Straub: (2022)

Theorem With the choice $a(n, k) = \binom{n}{k}^s$, the order of minimal recursion for $A^{(s)}(n)$ is at least $\lfloor \frac{s+1}{2} \rfloor$.

A bit of history:

1997 M. Skell showed the upper bound $\lfloor \frac{s+1}{2} \rfloor$
for the recursion for $A^{(s)}(n)$.

He uses estimates from Linear Algebra,
but not creative telescoping

(but his estimates also depend on representation
with $a(n, k) \geq \binom{n}{k}^s$!)

2007: Yuan, Lu & Schmidt uses an ingenious
congruence argument to give the lower bound
 ≥ 2 for every $A^{(s)}(n)$ with $s \geq 3$

Nowadays: for each fixed s , to show that
the operator is of minimal order, one
uses available algorithms from computer algebra.

By hand: up to $s \leq 20$.

Our result is for general s .

There is a different aspect of such recursions of order ≥ 2 : Apéry limits.

If one chooses two linearly independent solutions of such a recursion, $A(n)$ and $B(n)$ say,

then a question of interest is the limit

$$\lim_{n \rightarrow \infty} \frac{B(n)}{A(n)}$$

Apéry (1978) used the recursion

$$(n+1)^3 A(n+1) - (2n+1)(17n^2+17n+5)A(n) + n^3 A(n-1) = 0$$

for $A(n) = \sum_{k=0}^n \binom{n}{k}^2 \binom{n+k}{k}^2$ and another solution

$$B(n), B(0) = 20, B(1) = 6^5$$

to show that $\lim_{n \rightarrow \infty} \frac{B(n)}{A(n)} = \zeta(3)$.

In a recent paper, Chamberland & Straub conjectured that one can choose ^{indep.} solutions

$A_j^{(s)}(n)$ of the recursion for $A^{(s)}(n)$

such that $A_0^{(s)}(n) = A^{(s)}(n)$ and

$\lim_{n \rightarrow \infty} \frac{A_j^{(s)}(n)}{A_0^{(s)}(n)}$ is a ^{nonzero} rational multiple of n^{-2j} for $j = 0, 1, \dots, \lfloor \frac{s+1}{2} \rfloor - 1$.

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We confirm this observation by constructing explicitly such ~~as basis of~~ $\left[\frac{s+1}{2}\right]$ independent solutions (based on $a(n, k) \approx \binom{n}{k}^s$) and in fact use Asymptotic limits to prove that order in the main theorem is $\geq \left[\frac{s+1}{2}\right]$.

The idea of constructing the independent solutions is quite simple:

$$P(n, N) a(n, k) \approx B(n, k+1) - B(n, k)$$

can be used not only for integral k !

Using it for $a(n, k-t)$ where $t \in \mathbb{R}$ (say $|t| < 1$) we obtain that

$$A^{(s)}(n, t) \approx \left(\frac{1-t}{1+st}\right)^s \sum_{k=0}^n \binom{n}{k-t}^s$$

satisfies $P(n, N) A^{(s)}(n, t) \approx \underline{0} (t^s)$ as $t \rightarrow 0$

In other words, all coeffs in the t -expansion of $A^{(s)}(n, t)$ for t^j with $j \leq s$ satisfy the same recursion!

And $A^{(s)}(n, t)$ is an even function, one can check the non-zero coeffs of t^j are actually positive.

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$$\text{Finally, } \lim_{n \rightarrow \infty} \frac{A^{(s)}(n, t)}{A^{(0)}(n)} = \left(\frac{\pi t}{2n\pi t} \right)^s$$

$$= \left(\sum_{j=1}^{\infty} \left(2 - \frac{1}{2^{2j-2}} \right) \varphi(2j) t^{2j} \right)^s$$

$$= \left(\sum_{j=1}^{\infty} \left(\frac{1}{2^{2j-1}} - 1 \right) \frac{\beta_{2j}}{(2j)!} (2\pi t)^{2j} \right)^s$$

$$\in \mathbb{H} \otimes \mathbb{Q} \left[[\pi^2 t^2] \right]$$

The estimate for the order:

Assume the the solutions $A_j^{(s)}(n)$ so constructed are linearly dependent \Rightarrow lin dependent over \mathbb{Q} because $A_j^{(s)}(n) \in \mathbb{Q} \quad \forall j, n$

$$\text{write } 0 = \sum_{j=0}^{\lfloor \frac{s}{2} \rfloor - 1} \lambda_j A_j^{(s)}(n) \quad \lambda_j \in \mathbb{Q}$$

for all $n > n_0$.

Divide by $A_0^{(s)}(n)$ and take the limit as $n \rightarrow \infty$

$$0 = \sum_j \lambda_j \varphi_j \pi^{2j} \quad \text{where } \varphi_j \in \mathbb{Q}^{\times}$$

but the powers of π^2 are lin. independent over \mathbb{Q}

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