

# Differential Equation Invariance Axiomatization

André Platzer

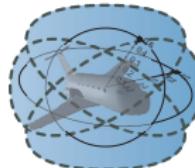
Karlsruhe Institute of Technology  
Department of Informatics

Computer Science Department  
Carnegie Mellon University

A. Platzer and Y.-K. Tan. J. ACM **67**(1), 2020.



Alexander von  
**HUMBOLDT**  
STIFTUNG



## 1 Differential Dynamic Logic

- Syntax
- Axiomatization
- Relative Completeness / ODE

## 2 Proofs for Differential Equations

- Differential Invariants / Cuts / Ghosts

## 3 Completeness for Differential Equation Invariants

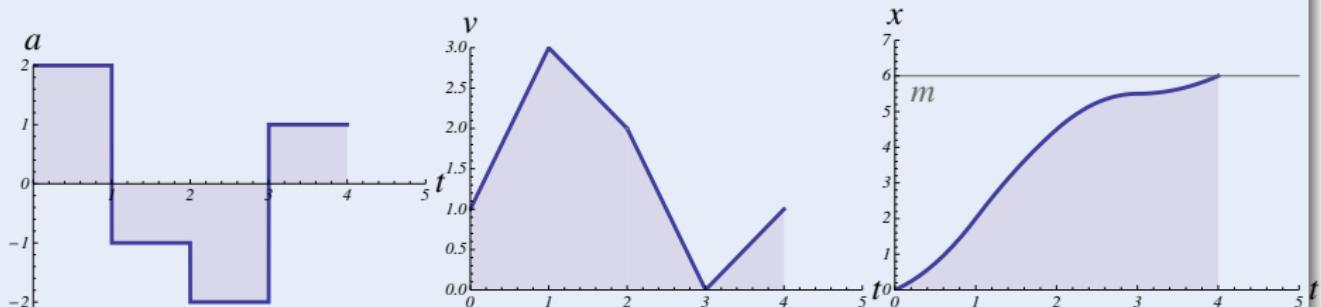
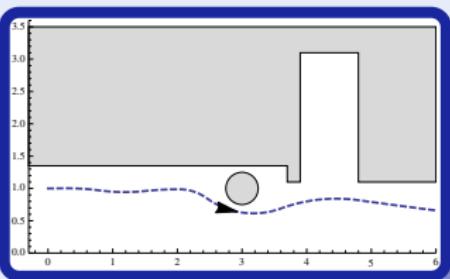
- Darboux are Differential Ghosts
- Derived Differential Radical Invariants
- Real Induction
- Derived Local Progress
- Completeness for Invariants
- Completeness for Noetherian Functions

## 4 Summary

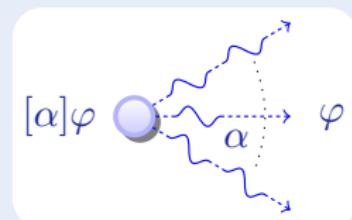
## Challenge (Hybrid Systems)

Fixed law describing state evolution with both

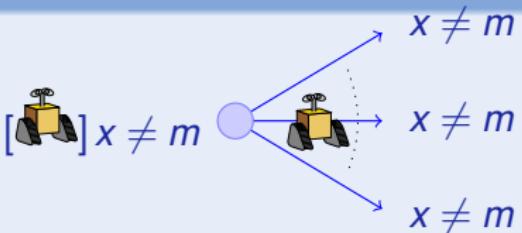
- Discrete dynamics (control decisions)
- Continuous dynamics (differential equations)



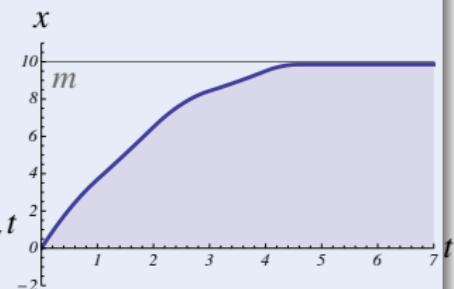
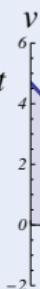
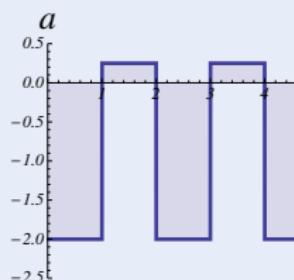
## Concept (Differential Dynamic Logic)



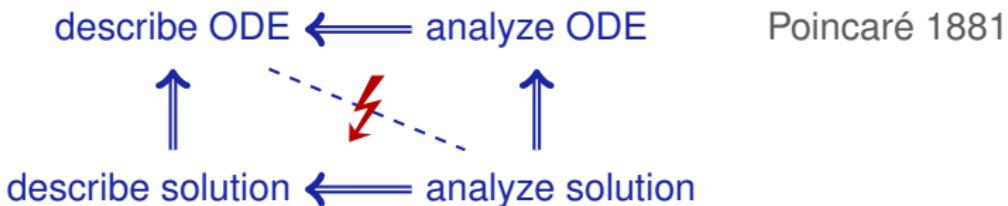
(JAR'08,LICS'12)



$$\underbrace{x \neq m \wedge b > 0}_{\text{init}} \rightarrow \left[ \left( \underbrace{\text{if}(\text{SB}(x, m)) a := -b}_{\text{all runs}} ; x' = v, v' = a \right)^* \right] \underbrace{x \neq m}_{\text{post}}$$



- Classical approach: ① Given ODE ② Solve ODE ③ Analyze solution
- Descriptive power of ODEs: ODE much easier than its solution
- ⚡ Analyzing ODEs via their solutions undoes their descriptive power!



- ① Now: Logical foundations of differential equation invariants
- ② Identify axioms for differential equations
- ③ Completeness for differential equation invariants
- ④ Uniformly substitutable axioms, not infinite axiom schemata
- ⑤ Decide invariance by proof

## 1 Differential Dynamic Logic

- Syntax
- Axiomatization
- Relative Completeness / ODE

## 2 Proofs for Differential Equations

- Differential Invariants / Cuts / Ghosts

## 3 Completeness for Differential Equation Invariants

- Darboux are Differential Ghosts
- Derived Differential Radical Invariants
- Real Induction
- Derived Local Progress
- Completeness for Invariants
- Completeness for Noetherian Functions

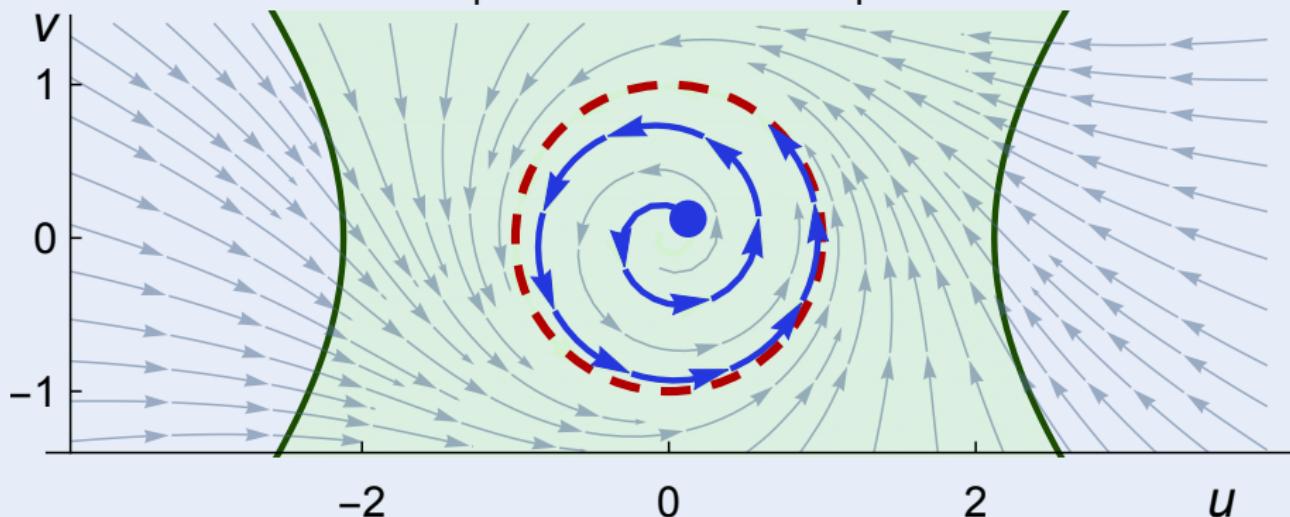
## 4 Summary

## Concept (Differential Dynamic Logic)

(JAR'08,LICS'12)

$$u^2 \leq v^2 + \frac{9}{2} \rightarrow [u' = -v + \frac{u}{4}(1-u^2-v^2), v' = u + \frac{v}{4}(1-u^2-v^2)] \quad u^2 \leq v^2 + \frac{9}{2}$$

$$u^2 + v^2 = 1 \rightarrow [u' = -v + \frac{u}{4}(1-u^2-v^2), v' = u + \frac{v}{4}(1-u^2-v^2)] \quad u^2 + v^2 = 1$$



Definition (Hybrid program  $\alpha$ )

$$\alpha, \beta ::= x := f(x) \mid ?Q \mid \textcolor{red}{x' = f(x) \& Q} \mid \alpha \cup \beta \mid \alpha ; \beta \mid \alpha^*$$

Definition (dL Formula  $P$ )

$$P, Q ::= e \geq \tilde{e} \mid \neg P \mid P \wedge Q \mid \forall x P \mid \exists x P \mid [\alpha]P \mid \langle \alpha \rangle P$$

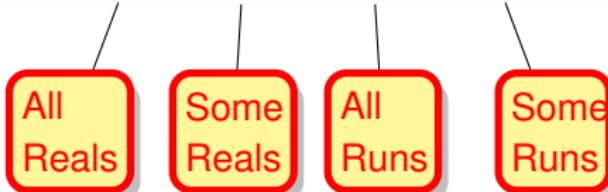


Definition (Hybrid program  $\alpha$ )

$$\alpha, \beta ::= x := f(x) \mid ?Q \mid \textcolor{red}{x' = f(x) \& Q} \mid \alpha \cup \beta \mid \alpha ; \beta \mid \alpha^*$$

Definition (dL Formula  $P$ )

$$P, Q ::= e \geq \tilde{e} \mid \neg P \mid P \wedge Q \mid \forall x P \mid \exists x P \mid [\alpha]P \mid \langle \alpha \rangle P$$



$$[:=] \quad [x := e]p(x) \leftrightarrow p(e)$$

equations of truth

$$[?] \quad [?Q]P \leftrightarrow (Q \rightarrow P)$$

$$['] \quad [x' = f(x)]p(x) \leftrightarrow \forall t \geq 0 [x := y(t)]p(x) \quad (y'(t) = f(y))$$

$$[\cup] \quad [\alpha \cup \beta]P \leftrightarrow [\alpha]P \wedge [\beta]P$$

$$[:] \quad [\alpha; \beta]P \leftrightarrow [\alpha][\beta]P$$

$$[*] \quad [\alpha^*]P \leftrightarrow P \wedge [\alpha][\alpha^*]P$$

$$\mathsf{K} \quad [\alpha](P \rightarrow Q) \rightarrow ([\alpha]P \rightarrow [\alpha]Q)$$

$$\mathsf{I} \quad [\alpha^*](P \rightarrow [\alpha]P) \rightarrow (P \rightarrow [\alpha^*]P)$$

$$\mathsf{C} \quad [\alpha^*]\forall v > 0 (P(v) \rightarrow \langle \alpha \rangle P(v-1)) \rightarrow \forall v (P(v) \rightarrow \langle \alpha^* \rangle \exists v \leq 0 P(v))$$

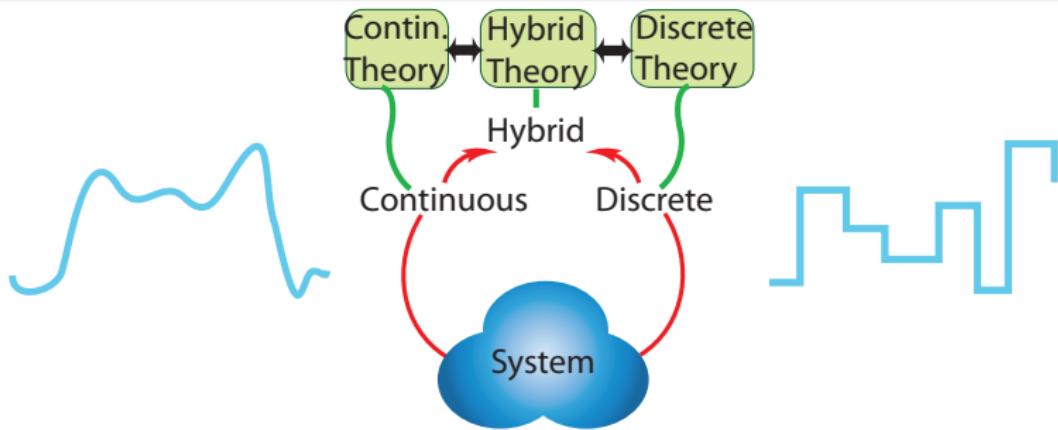
Theorem (Sound & Complete)

(JAR'08, LICS'12, JAR'17)

dL calculus is a sound & complete axiomatization of hybrid systems relative to either differential equations **or** relative to discrete dynamics.

Corollary (Complete Proof-theoretical Bridge)

proving continuous = proving hybrid = proving discrete



## 1 Differential Dynamic Logic

- Syntax
- Axiomatization
- Relative Completeness / ODE

## 2 Proofs for Differential Equations

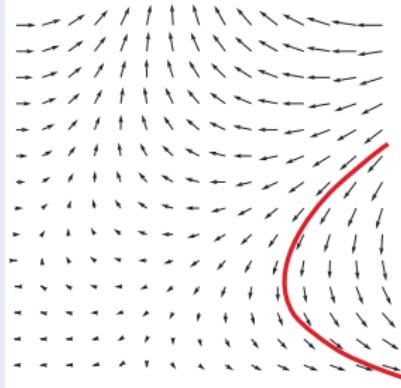
- Differential Invariants / Cuts / Ghosts

## 3 Completeness for Differential Equation Invariants

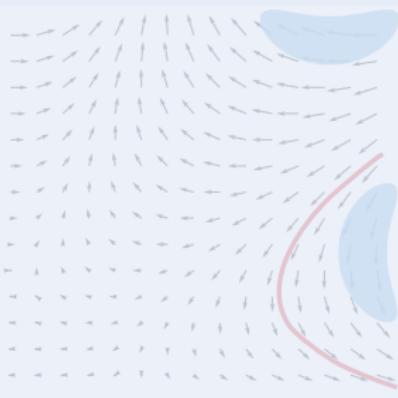
- Darboux are Differential Ghosts
- Derived Differential Radical Invariants
- Real Induction
- Derived Local Progress
- Completeness for Invariants
- Completeness for Noetherian Functions

## 4 Summary

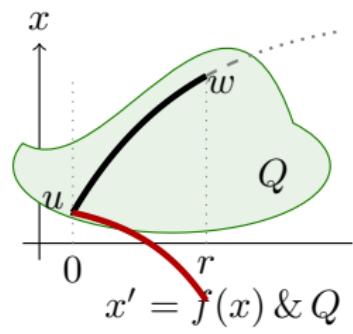
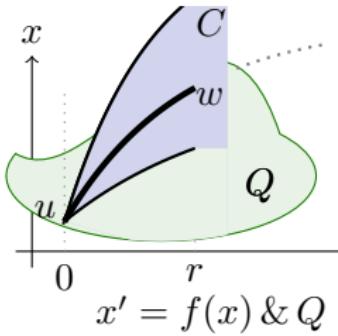
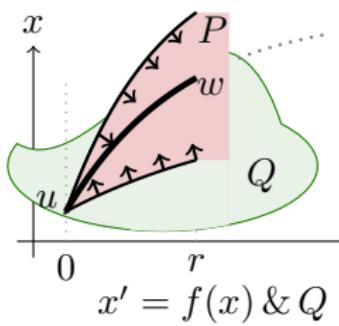
## Differential Invariant



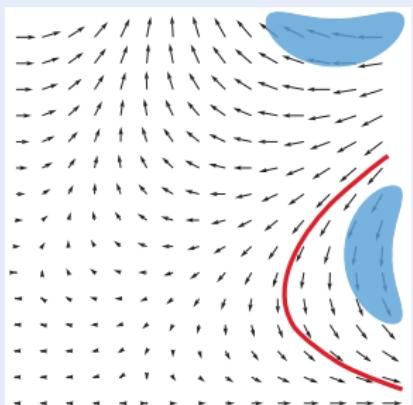
## Differential Cut



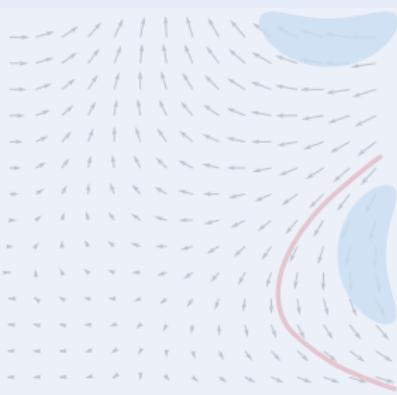
## Differential Ghost



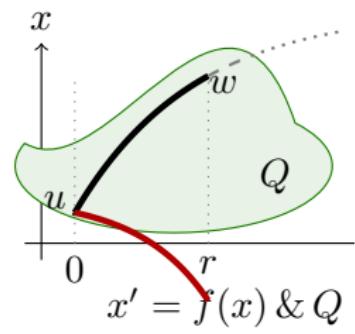
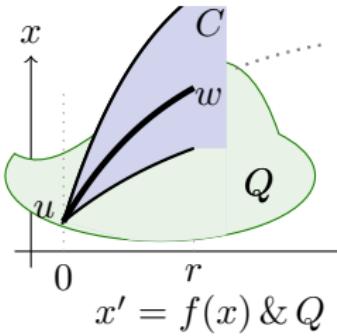
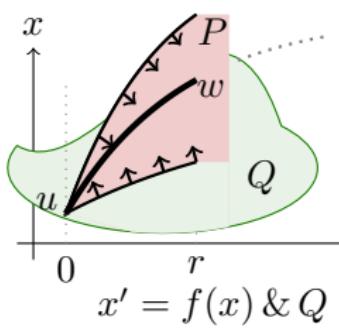
## Differential Invariant



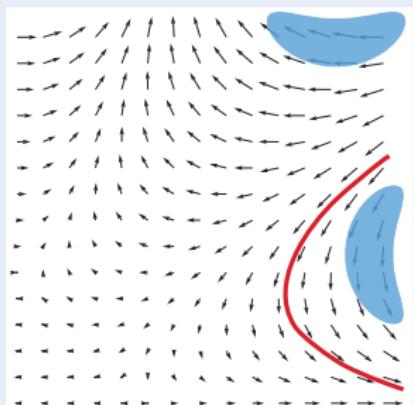
## Differential Cut



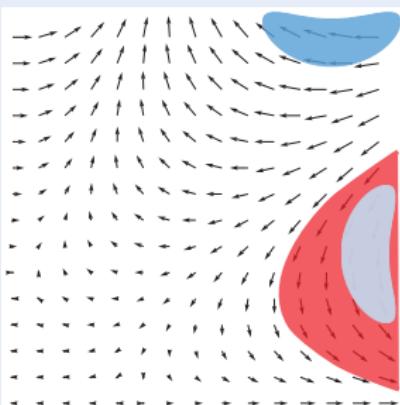
## Differential Ghost



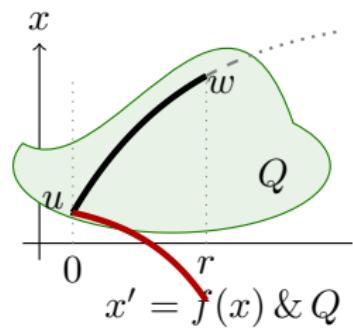
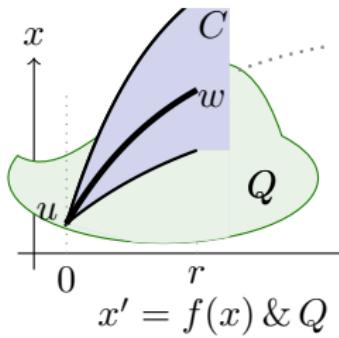
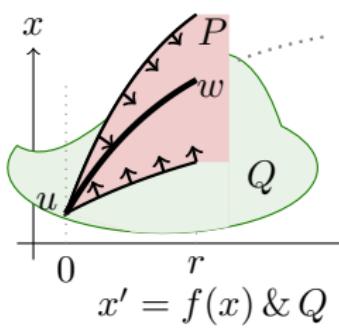
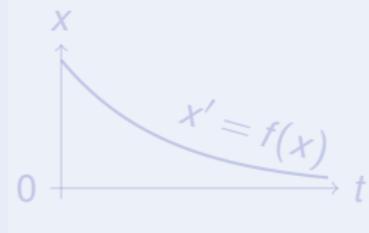
## Differential Invariant



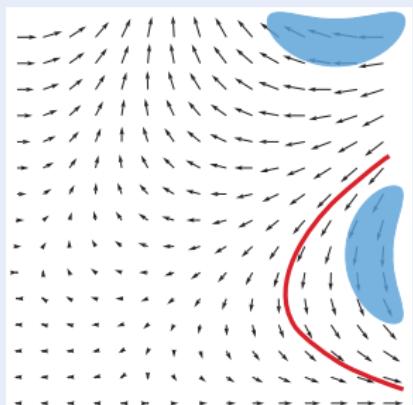
## Differential Cut



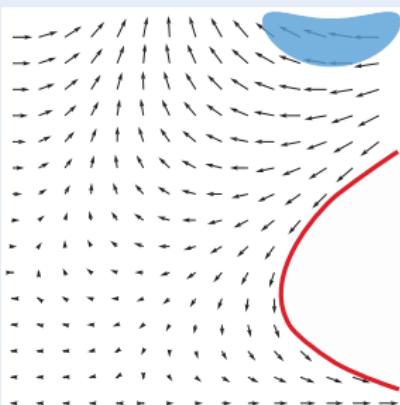
## Differential Ghost



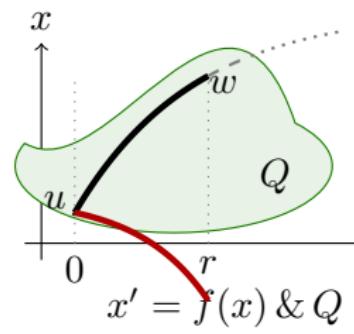
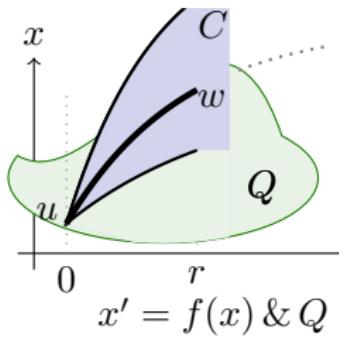
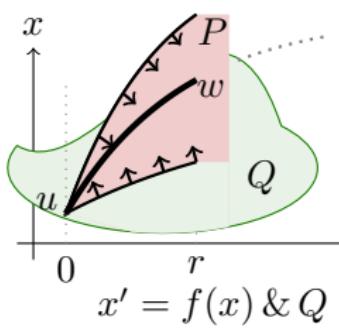
## Differential Invariant



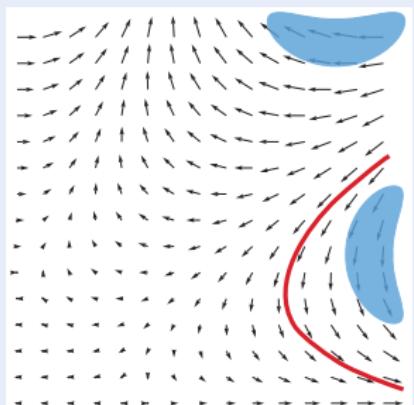
## Differential Cut



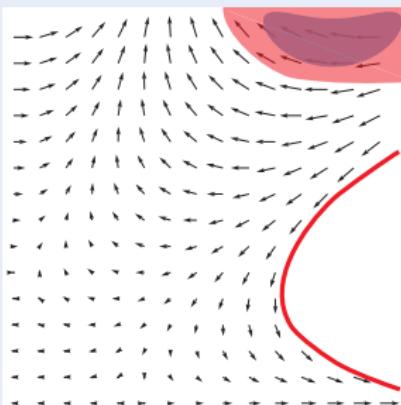
## Differential Ghost



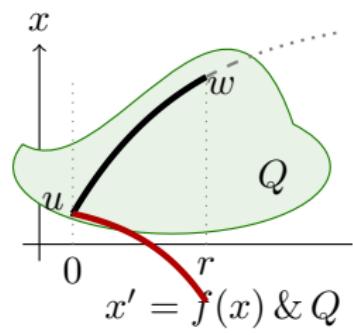
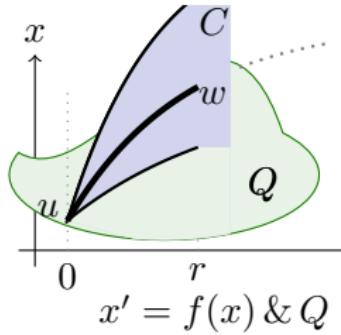
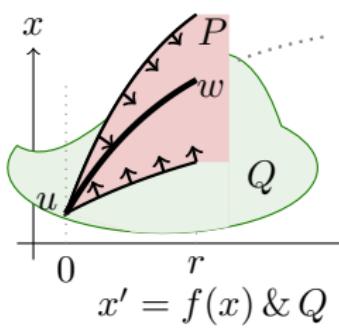
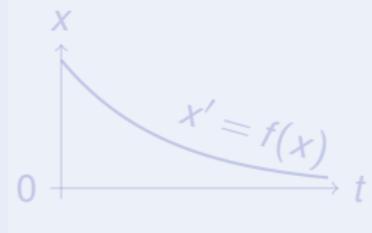
## Differential Invariant



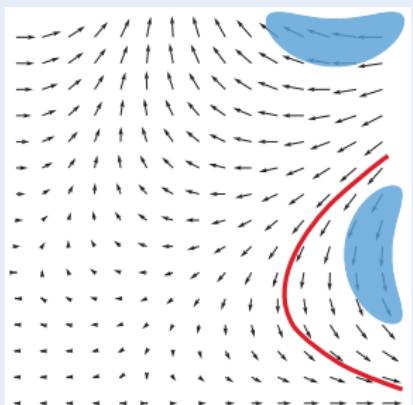
## Differential Cut



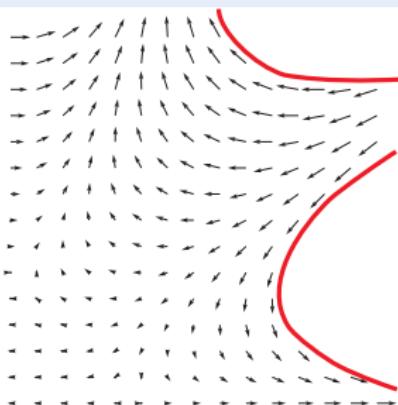
## Differential Ghost



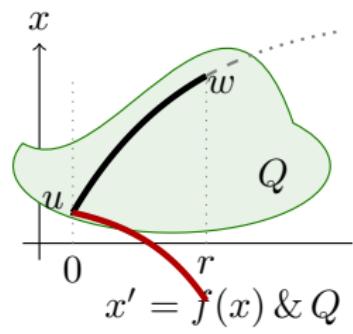
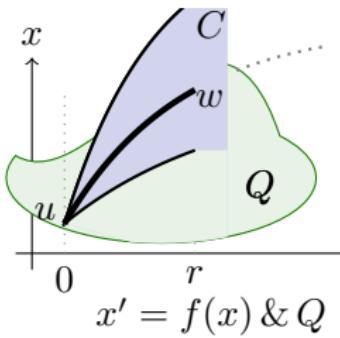
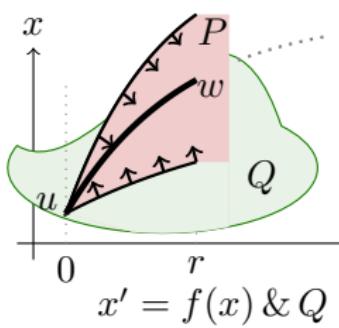
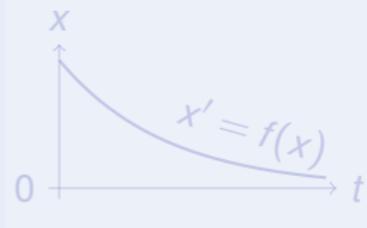
## Differential Invariant



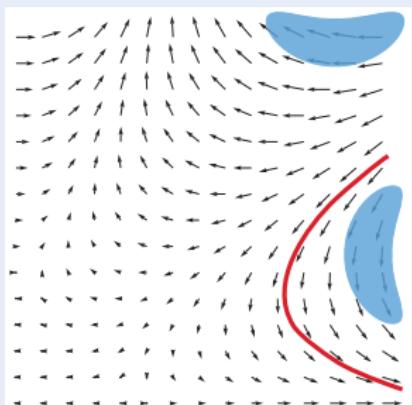
## Differential Cut



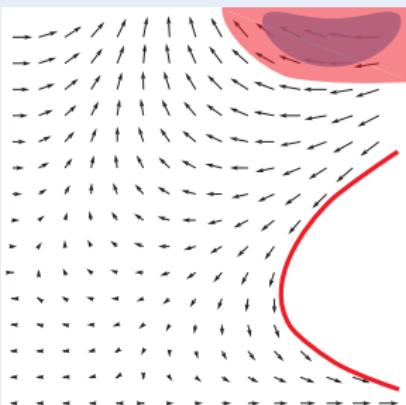
## Differential Ghost



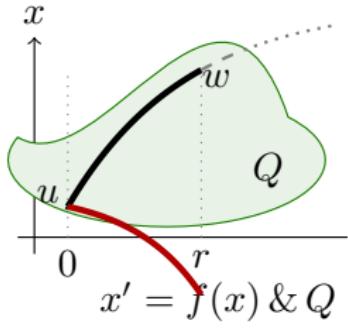
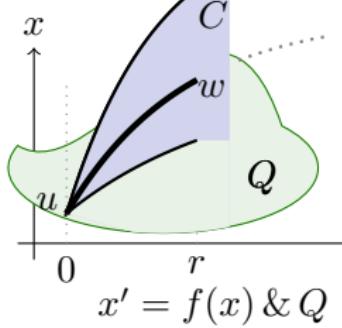
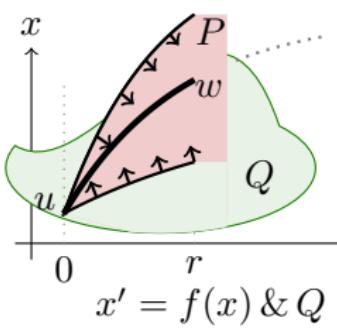
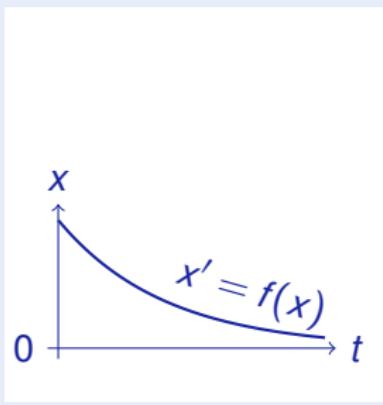
## Differential Invariant



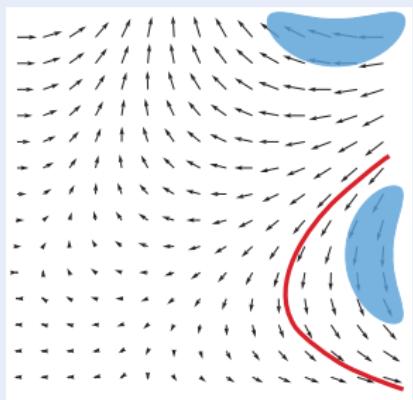
## Differential Cut



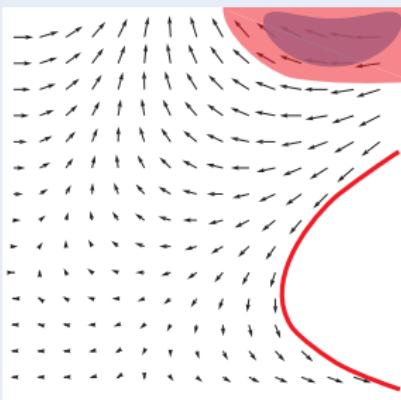
## Differential Ghost



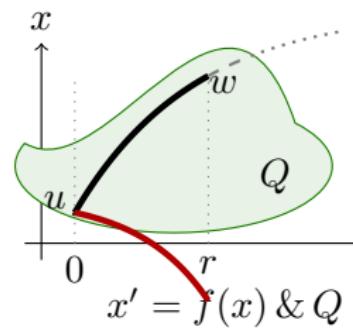
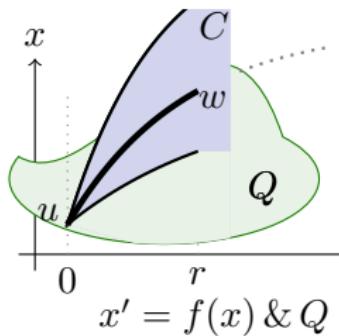
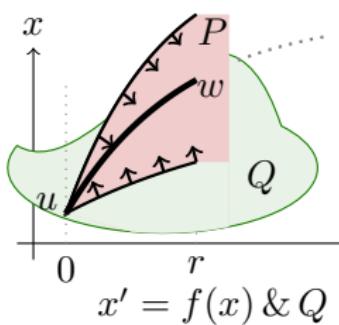
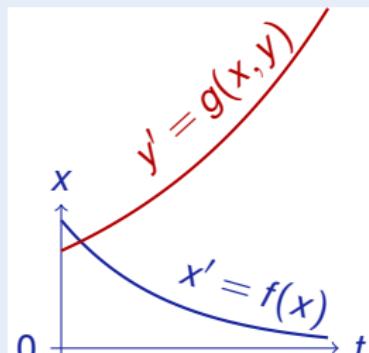
## Differential Invariant



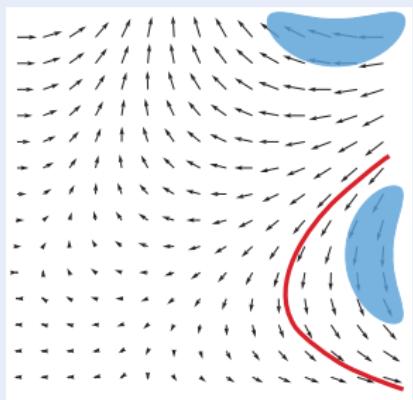
## Differential Cut



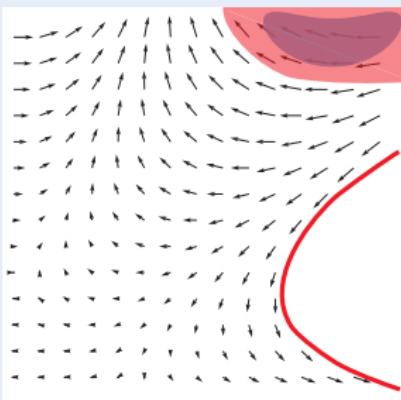
## Differential Ghost



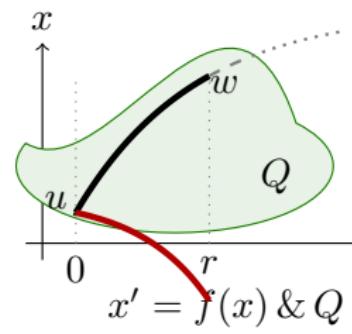
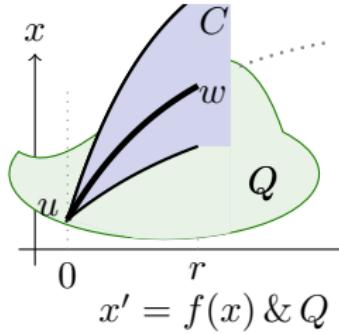
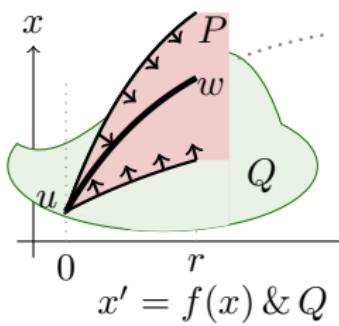
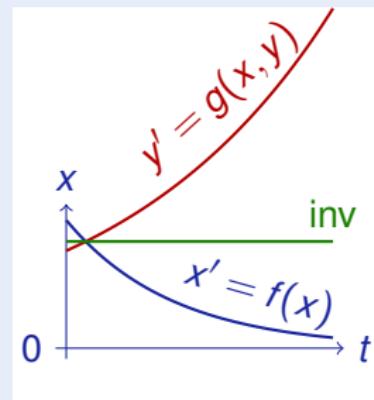
## Differential Invariant



## Differential Cut



## Differential Ghost



## Differential Invariant

$$\frac{Q \rightarrow [x' := f(x)](P)'}{P \rightarrow [x' = f(x) \& Q]P}$$

## Differential Cut

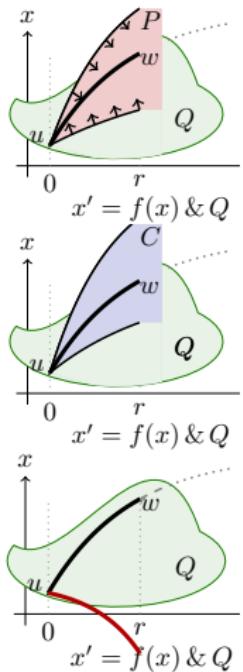
$$\frac{P \rightarrow [x' = f(x) \& Q]C \quad P \rightarrow [x' = f(x) \& Q \wedge C]P}{P \rightarrow [x' = f(x) \& Q]P}$$

## Differential Ghost

$$\frac{P \leftrightarrow \exists y G \quad G \rightarrow [x' = f(x), y' = g(x, y) \& Q]G}{P \rightarrow [x' = f(x) \& Q]P}$$

deductive power added DI  $\prec$  DI+DC  $\prec$  DI+DC+DG

$$v[[e]'] = \sum_x v(x') \frac{\partial [[e]]}{\partial x}(v)$$



## Differential Invariant

$$\frac{Q \rightarrow [x' := f(x)](P)'}{P \rightarrow [x' = f(x) \& Q]P}$$

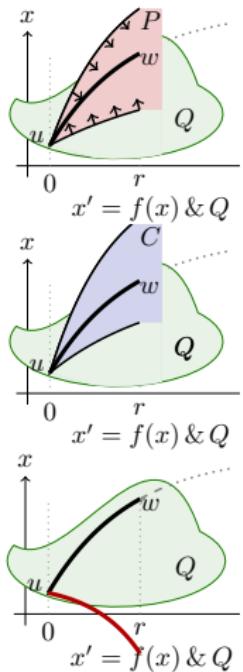
## Differential Cut

$$\frac{P \rightarrow [x' = f(x) \& Q]C \quad P \rightarrow [x' = f(x) \& Q \wedge C]P}{P \rightarrow [x' = f(x) \& Q]P}$$

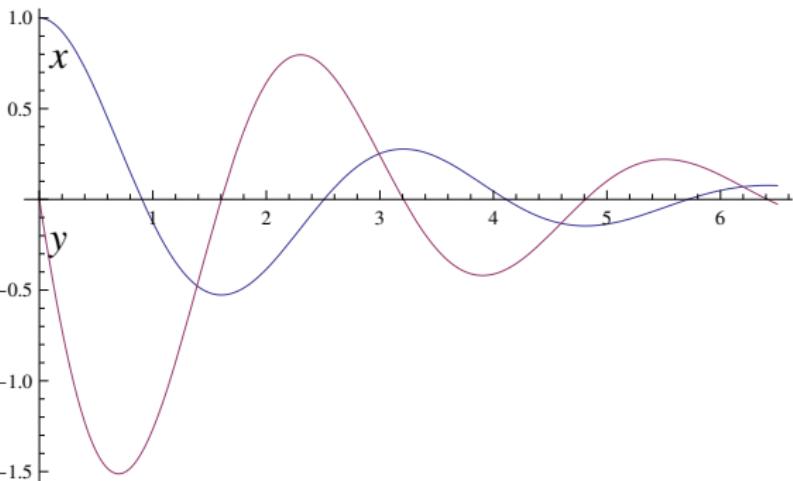
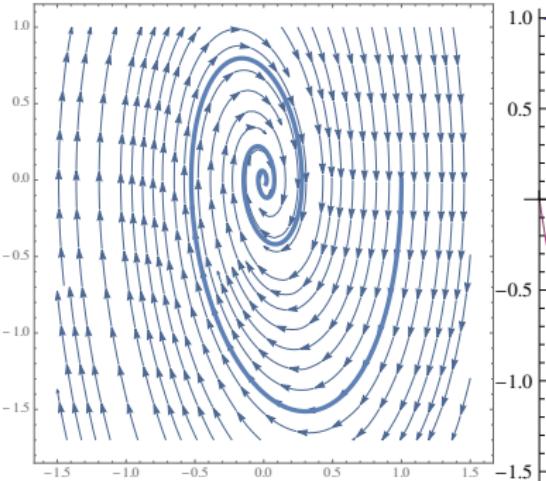
## Differential Ghost

$$\frac{P \leftrightarrow \exists y G \quad G \rightarrow [x' = f(x), y' = g(x, y) \& Q]G}{P \rightarrow [x' = f(x) \& Q]P}$$

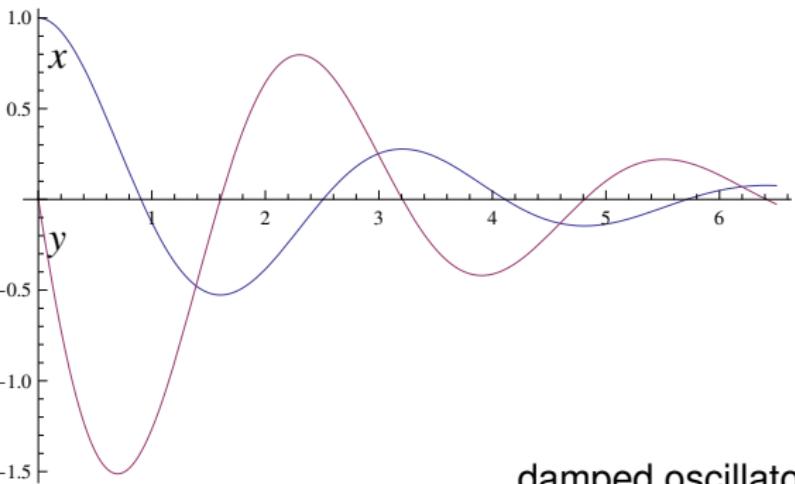
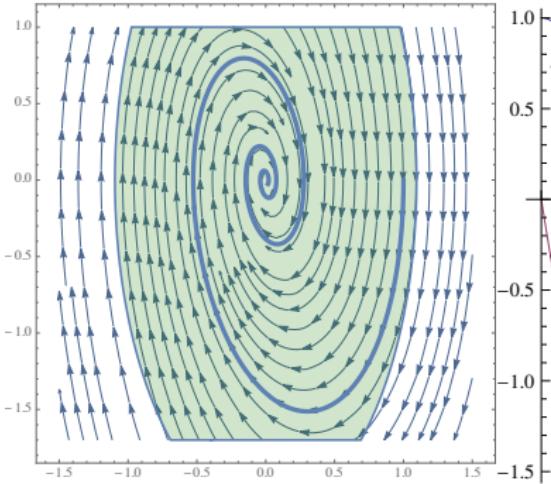
if  $g(x, y) = a(x)y + b(x)$ , so has long solution!



$$\omega^2 x^2 + y^2 \leq c^2 \rightarrow [x' = y, y' = -\omega^2 x - 2d\omega y \& \omega \geq 0 \wedge d \geq 0] \omega^2 x^2 + y^2 \leq c^2$$



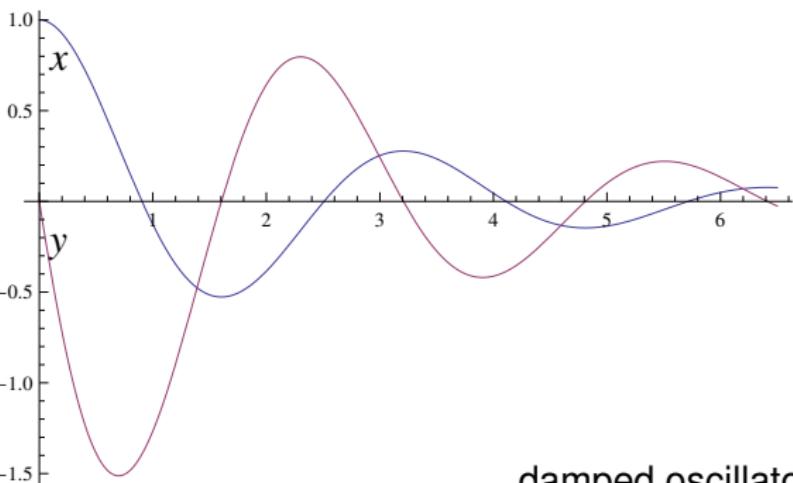
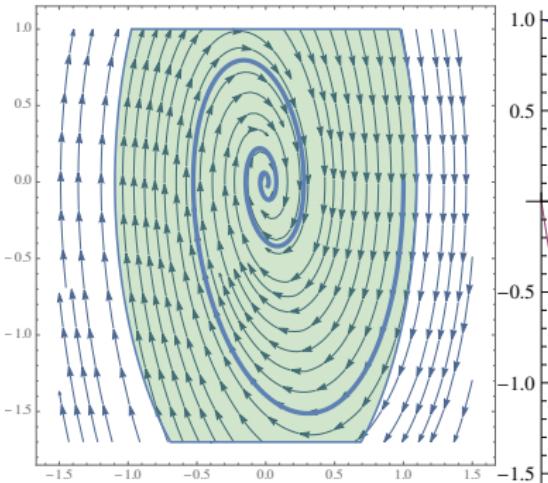
$$\omega^2 x^2 + y^2 \leq c^2 \rightarrow [x' = y, y' = -\omega^2 x - 2d\omega y \& \omega \geq 0 \wedge d \geq 0] \quad \omega^2 x^2 + y^2 \leq c^2$$



damped oscillator

$$\omega \geq 0 \wedge d \geq 0 \rightarrow [x' := y][y' := -\omega^2 x - 2d\omega y] 2\omega^2 x x' + 2y y' \leq 0$$

$$\omega^2 x^2 + y^2 \leq c^2 \rightarrow [x' = y, y' = -\omega^2 x - 2d\omega y \& \omega \geq 0 \wedge d \geq 0] \omega^2 x^2 + y^2 \leq c^2$$

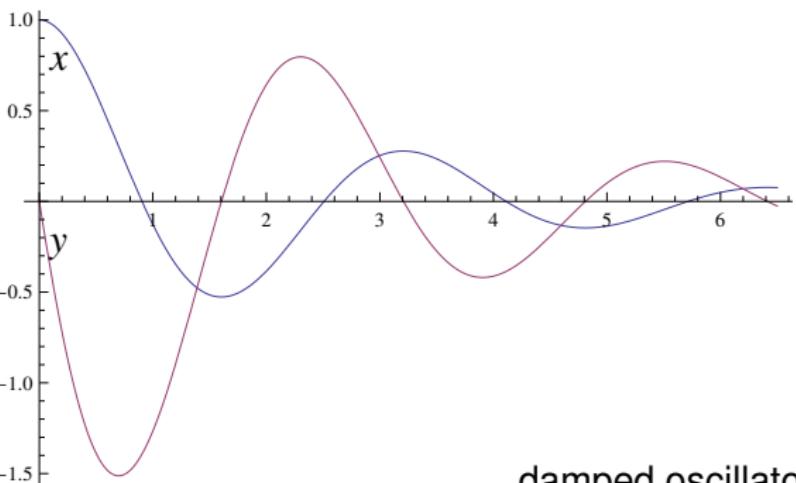
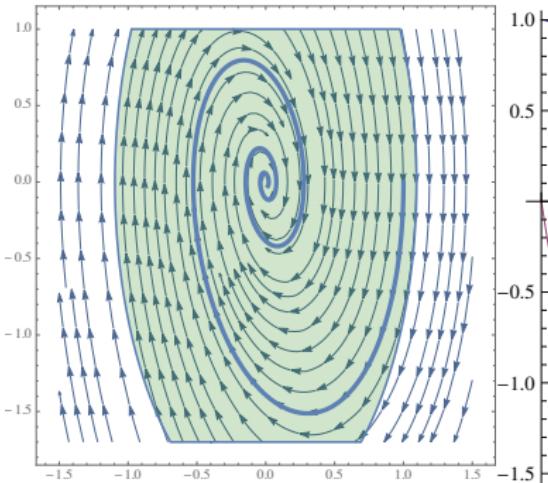


damped oscillator

$$\omega \geq 0 \wedge d \geq 0 \rightarrow 2\omega^2 xy + 2y(-\omega^2 x - 2d\omega y) \leq 0$$

$$\omega \geq 0 \wedge d \geq 0 \rightarrow [x' := y][y' := -\omega^2 x - 2d\omega y] 2\omega^2 x x' + 2y y' \leq 0$$

$$\omega^2 x^2 + y^2 \leq c^2 \rightarrow [x' = y, y' = -\omega^2 x - 2d\omega y \& \omega \geq 0 \wedge d \geq 0] \omega^2 x^2 + y^2 \leq c^2$$



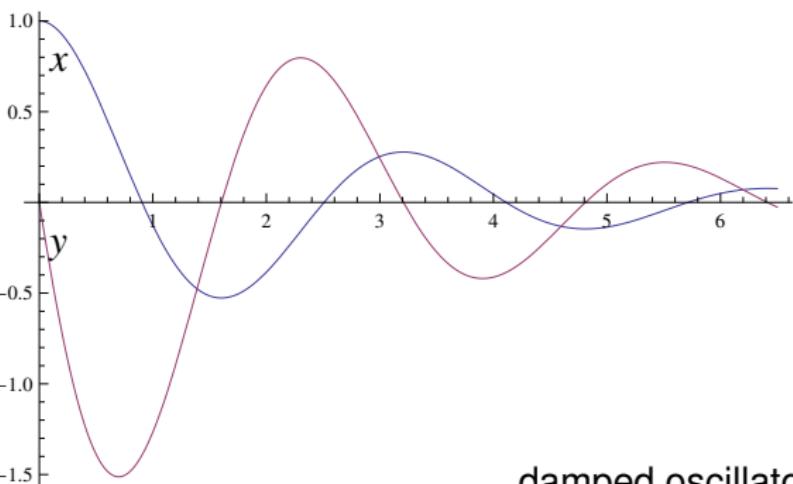
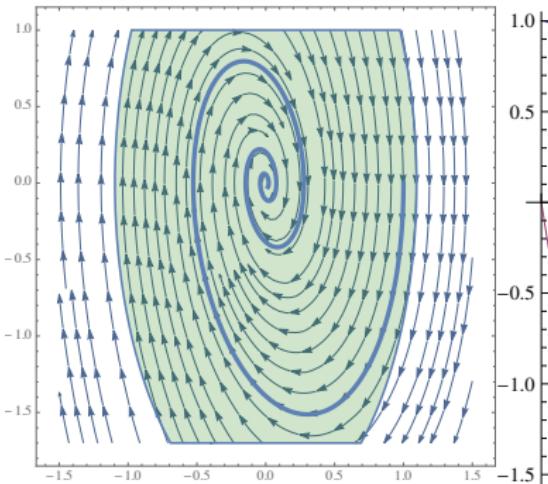
damped oscillator

\*

$$\omega \geq 0 \wedge d \geq 0 \rightarrow 2\omega^2 xy + 2y(-\omega^2 x - 2d\omega y) \leq 0$$

$$\omega \geq 0 \wedge d \geq 0 \rightarrow [x' := y][y' := -\omega^2 x - 2d\omega y] 2\omega^2 xx' + 2yy' \leq 0$$

$$\omega^2 x^2 + y^2 \leq c^2 \rightarrow [x' = y, y' = -\omega^2 x - 2d\omega y \& \omega \geq 0 \wedge d \geq 0] \omega^2 x^2 + y^2 \leq c^2$$



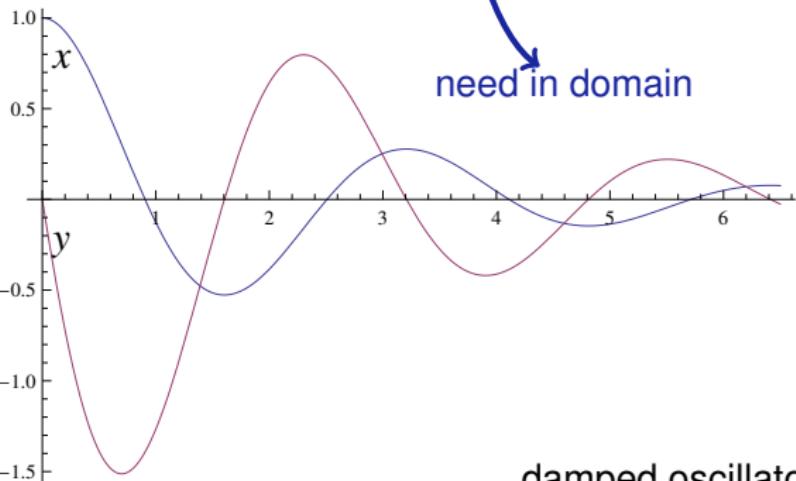
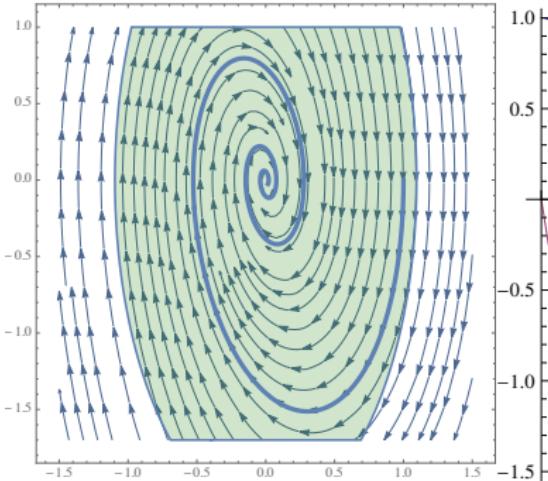
damped oscillator

\*

$$\omega \geq 0 \wedge d \geq 0 \rightarrow 2\omega^2 xy + 2y(-\omega^2 x - 2d\omega y) \leq 0$$

$$\omega \geq 0 \wedge d \geq 0 \rightarrow [x' := y][y' := -\omega^2 x - 2d\omega y] 2\omega^2 xx' + 2yy' \leq 0$$

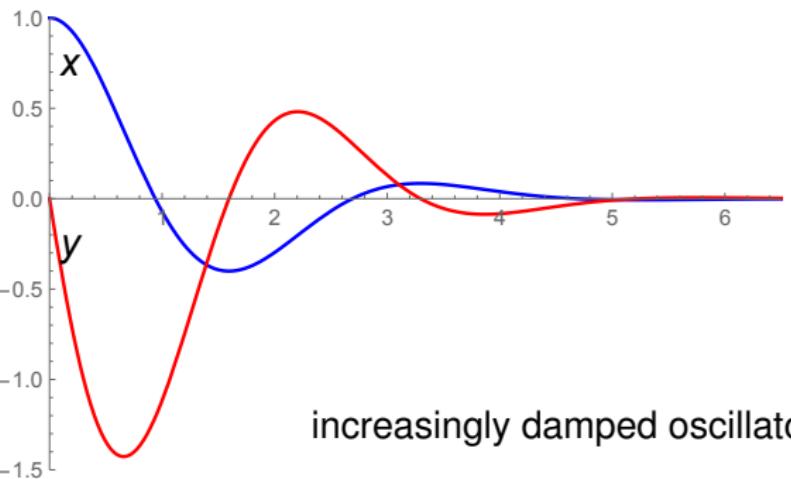
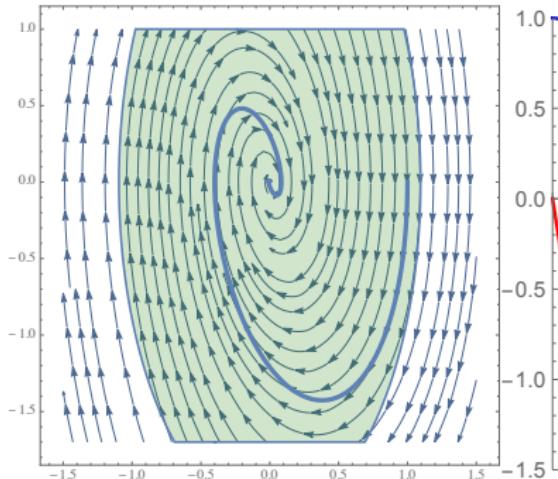
$$\omega^2 x^2 + y^2 \leq c^2 \rightarrow [x' = y, y' = -\omega^2 x - 2d\omega y \& \omega \geq 0 \wedge d \geq 0] \omega^2 x^2 + y^2 \leq c^2$$



---

$$\omega^2 x^2 + y^2 \leq c^2 \rightarrow [x' = y, y' = -\omega^2 x - 2d\omega y, d'=7 \text{ & } \omega \geq 0] \omega^2 x^2 + y^2 \leq c^2$$

$$\omega^2 x^2 + y^2 \leq c^2 \rightarrow [x' = y, y' = -\omega^2 x - 2d\omega y, d' = 7 \text{ & } \omega \geq 0] \quad \omega^2 x^2 + y^2 \leq c^2$$



$$\frac{\omega^2 x^2 + y^2 \leq c^2 \rightarrow [x' = y, y' = -\omega^2 x - 2d\omega y, d' = 7 \text{ & } \omega \geq 0 \wedge d \geq 0] \omega^2 x^2 + y^2 \leq c^2}{\omega^2 x^2 + y^2 \leq c^2 \rightarrow [x' = y, y' = -\omega^2 x - 2d\omega y, d' = 7 \text{ & } \omega \geq 0] \omega^2 x^2 + y^2 \leq c^2}$$

increasingly damped oscillator

$$\frac{\omega^2 x^2 + y^2 \leq c^2 \rightarrow [x' = y, y' = -\omega^2 x - 2d\omega y, d' = 7 \& \omega \geq 0 \wedge d \geq 0] \omega^2 x^2 + y^2 \leq c^2}{\omega^2 x^2 + y^2 \leq c^2 \rightarrow [x' = y, y' = -\omega^2 x - 2d\omega y, d' = 7 \& \omega \geq 0] \omega^2 x^2 + y^2 \leq c^2}$$

ask

$$d \geq 0 \rightarrow [x' = y, y' = -\omega^2 x - 2d\omega y, d' = 7 \& \omega \geq 0] d \geq 0$$

increasingly damped oscillator

$$\frac{\omega^2 x^2 + y^2 \leq c^2 \rightarrow [x' = y, y' = -\omega^2 x - 2d\omega y, d' = 7 \text{ & } \omega \geq 0 \wedge d \geq 0] \omega^2 x^2 + y^2 \leq c^2}{\omega^2 x^2 + y^2 \leq c^2 \rightarrow [x' = y, y' = -\omega^2 x - 2d\omega y, d' = 7 \text{ & } \omega \geq 0] \omega^2 x^2 + y^2 \leq c^2}$$

$$\omega \geq 0 \rightarrow [d' := 7] d' \geq 0$$

$$d \geq 0 \rightarrow [x' = y, y' = -\omega^2 x - 2d\omega y, d' = 7 \text{ & } \omega \geq 0] d \geq 0$$

increasingly damped oscillator

$$\frac{\omega^2 x^2 + y^2 \leq c^2 \rightarrow [x' = y, y' = -\omega^2 x - 2d\omega y, d' = 7 \text{ & } \omega \geq 0 \wedge d \geq 0] \omega^2 x^2 + y^2 \leq c^2}{\omega^2 x^2 + y^2 \leq c^2 \rightarrow [x' = y, y' = -\omega^2 x - 2d\omega y, d' = 7 \text{ & } \omega \geq 0] \omega^2 x^2 + y^2 \leq c^2}$$

---

$$\omega \geq 0 \rightarrow 7 \geq 0$$

---

$$\omega \geq 0 \rightarrow [d' := 7] d' \geq 0$$

---

$$d \geq 0 \rightarrow [x' = y, y' = -\omega^2 x - 2d\omega y, d' = 7 \text{ & } \omega \geq 0] d \geq 0$$

increasingly damped oscillator

$$\frac{\omega^2 x^2 + y^2 \leq c^2 \rightarrow [x' = y, y' = -\omega^2 x - 2d\omega y, d' = 7 \text{ & } \omega \geq 0 \wedge d \geq 0] \omega^2 x^2 + y^2 \leq c^2}{\omega^2 x^2 + y^2 \leq c^2 \rightarrow [x' = y, y' = -\omega^2 x - 2d\omega y, d' = 7 \text{ & } \omega \geq 0] \omega^2 x^2 + y^2 \leq c^2}$$

DC

\*

$$\omega \geq 0 \rightarrow 7 \geq 0$$

$$\omega \geq 0 \rightarrow [d' := 7] d' \geq 0$$

$$d \geq 0 \rightarrow [x' = y, y' = -\omega^2 x - 2d\omega y, d' = 7 \text{ & } \omega \geq 0] d \geq 0$$

increasingly damped oscillator

$$\omega \geq 0 \wedge d \geq 0 \rightarrow [x' := y][y' := -\omega^2 x - 2d\omega y] 2\omega^2 x x' + 2y y' \leq 0$$

$$\omega^2 x^2 + y^2 \leq c^2 \rightarrow [x' = y, y' = -\omega^2 x - 2d\omega y, d' = 7 \text{ & } \omega \geq 0 \wedge d \geq 0] \omega^2 x^2 + y^2 \leq c^2$$

$$\omega^2 x^2 + y^2 \leq c^2 \rightarrow [x' = y, y' = -\omega^2 x - 2d\omega y, d' = 7 \text{ & } \omega \geq 0] \omega^2 x^2 + y^2 \leq c^2$$

\*

$$\omega \geq 0 \rightarrow 7 \geq 0$$

$$\omega \geq 0 \rightarrow [d' := 7] d' \geq 0$$

$$d \geq 0 \rightarrow [x' = y, y' = -\omega^2 x - 2d\omega y, d' = 7 \text{ & } \omega \geq 0] d \geq 0$$

increasingly damped oscillator

$$\omega \geq 0 \wedge d \geq 0 \rightarrow 2\omega^2 xy + 2y(-\omega^2 x - 2d\omega y) \leq 0$$

$$\omega \geq 0 \wedge d \geq 0 \rightarrow [x' := y][y' := -\omega^2 x - 2d\omega y] 2\omega^2 x x' + 2y y' \leq 0$$

$$\omega^2 x^2 + y^2 \leq c^2 \rightarrow [x' = y, y' = -\omega^2 x - 2d\omega y, d' = 7 \text{ & } \omega \geq 0 \wedge d \geq 0] \omega^2 x^2 + y^2 \leq c^2$$

$$\omega^2 x^2 + y^2 \leq c^2 \rightarrow [x' = y, y' = -\omega^2 x - 2d\omega y, d' = 7 \text{ & } \omega \geq 0] \omega^2 x^2 + y^2 \leq c^2$$

\*

$$\omega \geq 0 \rightarrow 7 \geq 0$$

$$\omega \geq 0 \rightarrow [d' := 7] d' \geq 0$$

$$d \geq 0 \rightarrow [x' = y, y' = -\omega^2 x - 2d\omega y, d' = 7 \text{ & } \omega \geq 0] d \geq 0$$

increasingly damped oscillator

\*

$$\omega \geq 0 \wedge d \geq 0 \rightarrow 2\omega^2 xy + 2y(-\omega^2 x - 2d\omega y) \leq 0$$

$$\omega \geq 0 \wedge d \geq 0 \rightarrow [x' := y][y' := -\omega^2 x - 2d\omega y] 2\omega^2 xx' + 2yy' \leq 0$$

$$\omega^2 x^2 + y^2 \leq c^2 \rightarrow [x' = y, y' = -\omega^2 x - 2d\omega y, d' = 7 \text{ & } \omega \geq 0 \wedge d \geq 0] \omega^2 x^2 + y^2 \leq c^2$$

$$\omega^2 x^2 + y^2 \leq c^2 \rightarrow [x' = y, y' = -\omega^2 x - 2d\omega y, d' = 7 \text{ & } \omega \geq 0] \omega^2 x^2 + y^2 \leq c^2$$

\*

$$\omega \geq 0 \rightarrow 7 \geq 0$$

$$\omega \geq 0 \rightarrow [d' := 7] d' \geq 0$$

$$d \geq 0 \rightarrow [x' = y, y' = -\omega^2 x - 2d\omega y, d' = 7 \text{ & } \omega \geq 0] d \geq 0$$

increasingly damped oscillator

\*

$$\omega \geq 0 \wedge d \geq 0 \rightarrow 2\omega^2 xy + 2y(-\omega^2 x - 2d\omega y) \leq 0$$

$$\omega \geq 0 \wedge d \geq 0 \rightarrow [x' := y][y' := -\omega^2 x - 2d\omega y] 2\omega^2 xx' + 2yy' \leq 0$$

$$\omega^2 x^2 + y^2 \leq c^2 \rightarrow [x' = y, y' = -\omega^2 x - 2d\omega y, d' = 7 \text{ & } \omega \geq 0 \wedge d \geq 0] \omega^2 x^2 + y^2 \leq c^2$$

$$\omega^2 x^2 + y^2 \leq c^2 \rightarrow [x' = y, y' = -\omega^2 x - 2d\omega y, d' = 7 \text{ & } \omega \geq 0] \omega^2 x^2 + y^2 \leq c^2$$

init

\*

$$\omega \geq 0 \rightarrow 7 \geq 0$$

$$\omega \geq 0 \rightarrow [d' := 7] d' \geq 0$$

$$d \geq 0 \rightarrow [x' = y, y' = -\omega^2 x - 2d\omega y, d' = 7 \text{ & } \omega \geq 0] d \geq 0$$

Could repeatedly diffcut in formulas to help the proof

- 1 Differential Dynamic Logic
  - Syntax
  - Axiomatization
  - Relative Completeness / ODE
- 2 Proofs for Differential Equations
  - Differential Invariants / Cuts / Ghosts
- 3 Completeness for Differential Equation Invariants
  - Darboux are Differential Ghosts
  - Derived Differential Radical Invariants
  - Real Induction
  - Derived Local Progress
  - Completeness for Invariants
  - Completeness for Noetherian Functions
- 4 Summary

## Theorem (Algebraic Completeness)

(LICS'18,JACM'20)

dL calculus is a sound & complete axiomatization of algebraic invariants of polynomial differential equations. They are decidable by DI,DC,DG in dL.

## Theorem (Semialgebraic Completeness)

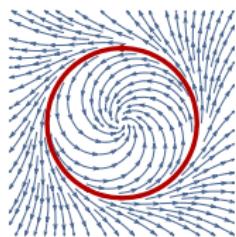
(LICS'18,JACM'20)

dL calculus with RI is a sound & complete axiomatization of semialgebraic invariants of differential equations. They are decidable in dL.

Darboux equalities are DG

Gaston Darboux 1878

$$\frac{Q \rightarrow e' = ge}{e = 0 \rightarrow [x' = f(x) \& Q]e = 0} \quad (g \in \mathbb{R}[x])$$

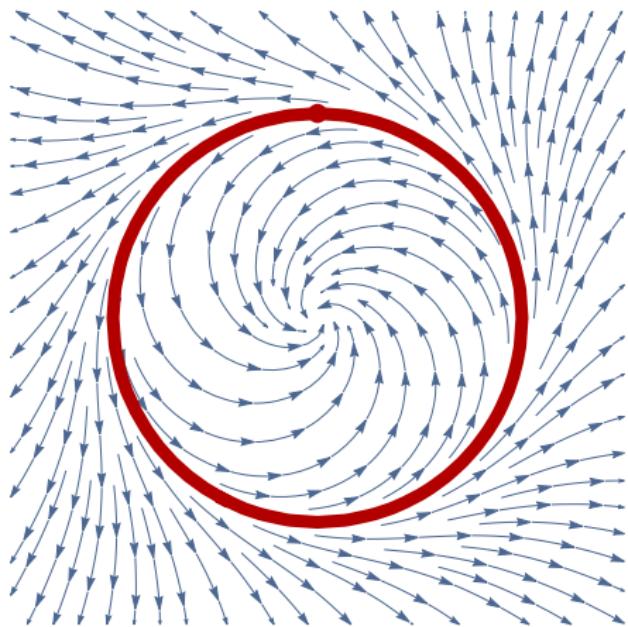
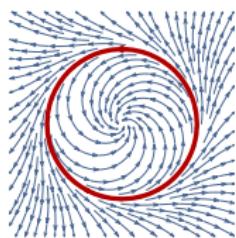


Definable  $e'$  for Lie-derivative w.r.t. ODE

Darboux equalities are DG

$$\frac{Q \rightarrow e' = ge}{e = 0 \rightarrow [x' = f(x) \& Q]e = 0} \quad (g \in \mathbb{R}[x])$$

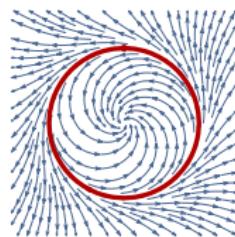
Gaston Darboux 1878



$$\frac{2uu' + 2vv' = 2(u^2 + v^2)(u^2 + v^2 - 1)}{\therefore \begin{aligned} u' &= -v - u + u^3 + uv^2 \\ v' &= u - v + u^2v + v^3 \end{aligned} \quad u^2 + v^2 - 1 = 0}$$

Darboux equalities are DG

$$\frac{Q \rightarrow e' = ge \quad (g \in \mathbb{R}[x])}{e = 0 \rightarrow [x' = f(x) \& Q]e = 0}$$



Proof Idea.

- ① DG counterweight  $y' = -gy$  to reduce  $e = 0$  to  $ey = 0 \wedge y \neq 0$ .
- ② DG counter-counterweight  $z' = gz$  to reduce  $y \neq 0$  to  $yz = 1$ .
- ③  $ey = 0$  and  $yz = 1$  are now differential invariants by construction.

□

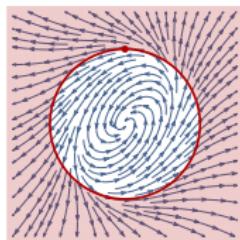
Derive

$$[x' = f(x) \& Q](e)' = ge \rightarrow (e = 0 \rightarrow [x' = f(x) \& Q]e = 0)$$

Darboux **inequalities** are DG

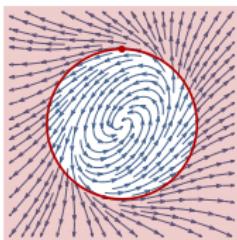
Thomas Grönwall 1919

$$\frac{Q \rightarrow e' \geq g e}{e \succcurlyeq 0 \rightarrow [x' = f(x) \& Q] e \succcurlyeq 0} \quad (g \in \mathbb{R}[x])$$



Darboux inequalities are DG

$$\frac{Q \rightarrow e' \geq ge \quad (g \in \mathbb{R}[x])}{e \succcurlyeq 0 \rightarrow [x' = f(x) \& Q] e \succcurlyeq 0}$$



Proof Idea.

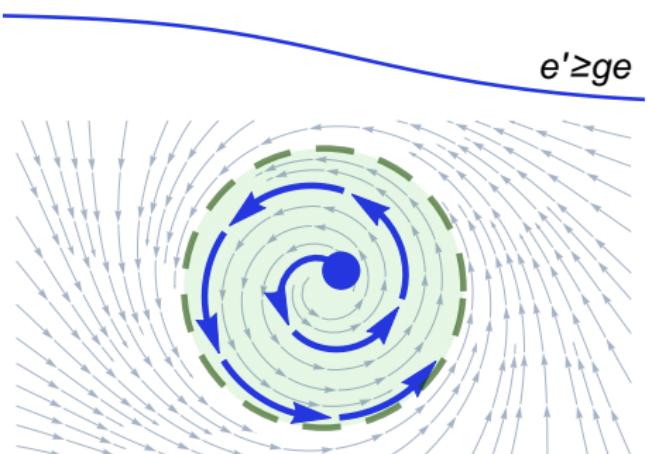
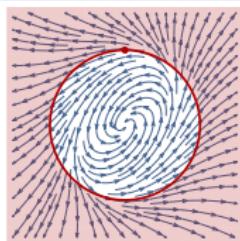
- ① DG counterweight  $y' = -gy$  to reduce  $e \succcurlyeq 0$  to  $ey \succcurlyeq 0 \wedge y > 0$ .
- ② DG counter-counterweight  $z' = \frac{g}{2}z$  to reduce  $y > 0$  to  $yz^2 = 1$ .
- ③  $yz^2 = 1$  and (after DC with  $y > 0$ )  $ey \succcurlyeq 0$  are differential invariants by construction as  $(ey)' = e'y - gye \geq 0$  from premise since  $y > 0$ . □

Derive

$$[x' = f(x) \& Q](e)' \geq ge \rightarrow (e \succcurlyeq 0 \rightarrow [x' = f(x) \& Q] e \succcurlyeq 0)$$

Darboux **inequalities** are DG

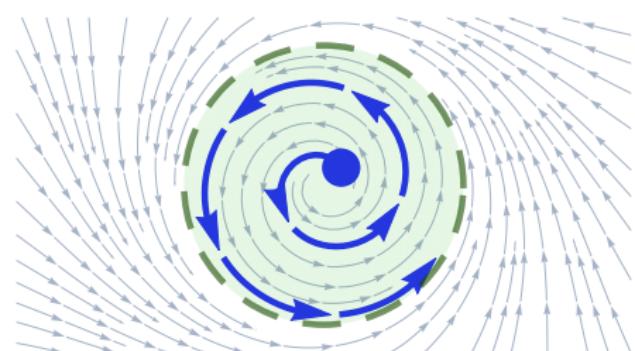
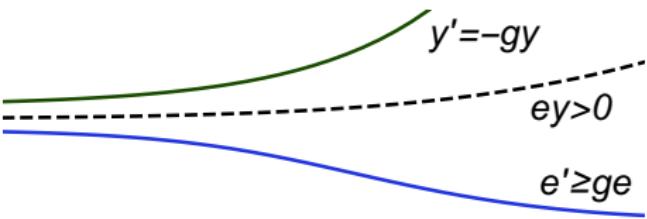
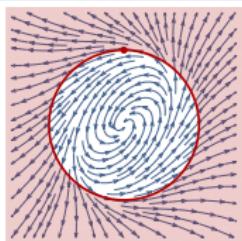
$$\frac{Q \rightarrow e' \geq ge}{e \succcurlyeq 0 \rightarrow [x' = f(x) \& Q] e \succcurlyeq 0} \quad (g \in \mathbb{R}[x])$$



$$\frac{(1-u^2-v^2)' \geq -\frac{1}{2}(u^2+v^2)(1-u^2-v^2)}{\dots \rightarrow [u' = -v + \frac{u}{4}(1-u^2-v^2) \\ v' = u + \frac{v}{4}(1-u^2-v^2) \\ ] 1-u^2-v^2 > 0}$$

Darboux **inequalities** are DG

$$\frac{Q \rightarrow e' \geq ge}{e \succcurlyeq 0 \rightarrow [x' = f(x) \& Q] e \succcurlyeq 0} \quad (g \in \mathbb{R}[x])$$

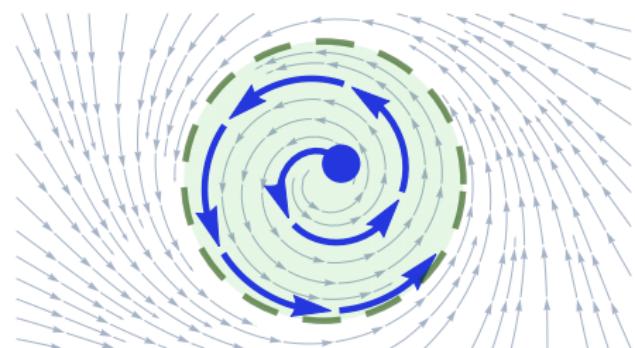
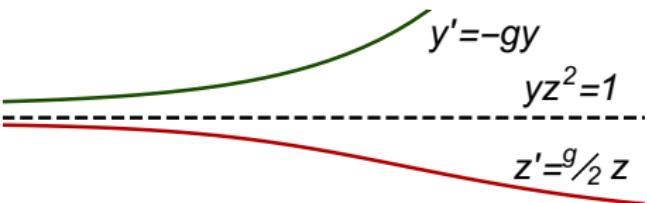
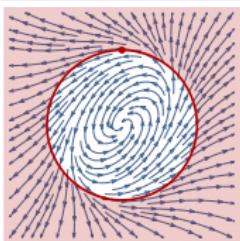


$$\begin{aligned} (1-u^2-v^2)' &\geq -\frac{1}{2}(u^2+v^2)(1-u^2-v^2) \\ \dots \rightarrow [u' &= -v + \frac{u}{4}(1-u^2-v^2) \\ v' &= u + \frac{v}{4}(1-u^2-v^2) \\ y' &= \frac{1}{2}(u^2+v^2)y \\ ] \quad 1-u^2-v^2 &> 0 \end{aligned}$$

$$(1-u^2-v^2)y > 0$$

Darboux **inequalities** are DG

$$\frac{Q \rightarrow e' \geq g e \quad (g \in \mathbb{R}[x])}{e \succcurlyeq 0 \rightarrow [x' = f(x) \& Q] e \succcurlyeq 0}$$



$$(1-u^2-v^2)' \geq -\frac{1}{2}(u^2+v^2)(1-u^2-v^2)$$

$$\dots \rightarrow [u' = -v + \frac{u}{4}(1-u^2-v^2)$$

$$v' = u + \frac{v}{4}(1-u^2-v^2)$$

$$y' = \frac{1}{2}(u^2+v^2)y$$

$$z' = -\frac{1}{4}(u^2+v^2)z$$

$$] 1-u^2-v^2 > 0$$

$$(1-u^2-v^2)y > 0$$

$$yz^2 = 1$$

$$\begin{array}{c}
 * \\
 \hline
 \text{R} \quad Q \rightarrow (-gy)z^2 + y(2z(\frac{g}{2}z)) = 0 \\
 \hline
 \text{dI} \quad yz^2 = 1 \rightarrow [x' = f(x), y' = -gy, z' = \frac{g}{2}z \& Q] yz^2 = 1 \\
 \hline
 \text{M,} \exists \text{R} \quad y > 0 \rightarrow \exists z [x' = f(x), y' = -gy, z' = \frac{g}{2}z \& Q] y > 0 \\
 \hline
 \text{dG} \quad y > 0 \rightarrow [x' = f(x), y' = -gy \& Q] y > 0 \\
 \end{array}$$

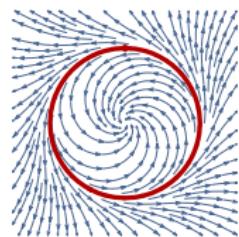
\*

$$\begin{array}{c}
 Q \rightarrow e' \geq ge \quad \overline{\text{R} \quad e' \geq ge, y > 0 \rightarrow e'y - gy \geq 0} \\
 \hline
 \text{cut} \quad Q, y > 0 \rightarrow e'y - gy \geq 0 \\
 \hline
 \text{dI} \quad e \succcurlyeq 0, y > 0 \rightarrow [x' = f(x), y' = -gy \& Q \wedge y > 0] ey \succcurlyeq 0 \triangleright \\
 \hline
 \text{dC} \quad e \succcurlyeq 0, y > 0 \rightarrow [x' = f(x), y' = -gy \& Q] (y > 0 \wedge ey \succcurlyeq 0) \\
 \hline
 \text{M,} \exists \text{R} \quad e \succcurlyeq 0 \rightarrow \exists y [x' = f(x), y' = -gy \& Q] e \succcurlyeq 0 \\
 \hline
 \text{dG} \quad e \succcurlyeq 0 \rightarrow [x' = f(x) \& Q] e \succcurlyeq 0
 \end{array}$$

P.S.  $z' = \frac{g}{2}z$  superfluous for open inequalities  $e > 0$  and  $e \neq 0$ .

Vectorial Darboux are DG

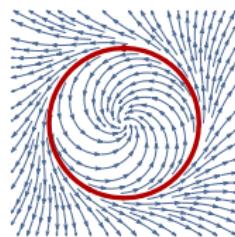
$$\frac{Q \rightarrow \mathbf{e}' = G\mathbf{e} \quad (G \in \mathbb{R}[x]^{n \times n})}{\mathbf{e} = 0 \rightarrow [x' = f(x) \& Q] \mathbf{e} = 0}$$



Definable  $\mathbf{e}'$  for component-wise Lie-derivative w.r.t. ODE

Vectorial Darboux are DG

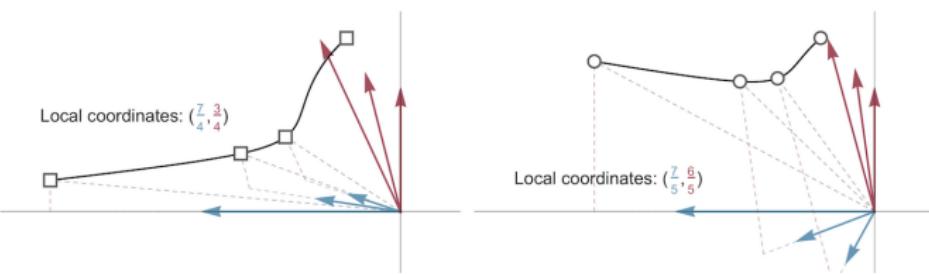
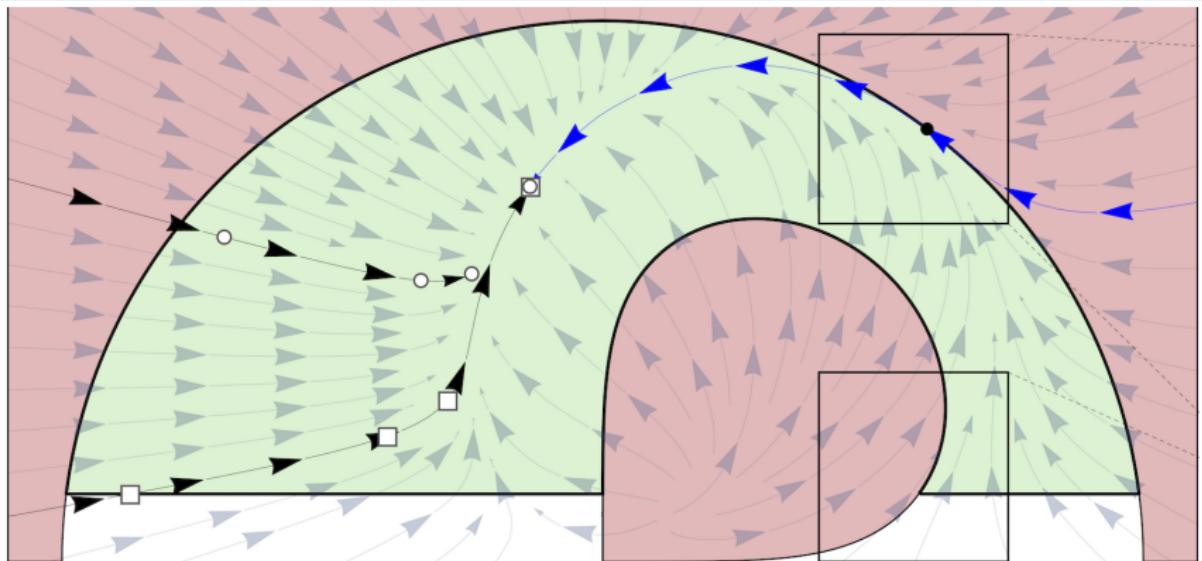
$$\frac{Q \rightarrow \mathbf{e}' = G\mathbf{e} \quad (G \in \mathbb{R}[x]^{n \times n})}{\mathbf{e} = 0 \rightarrow [x' = f(x) \& Q]\mathbf{e} = 0}$$



Proof Idea.

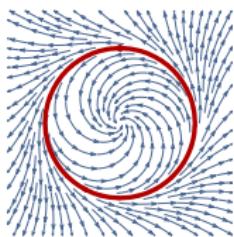
- ① DG counterweight  $\mathbf{y}' = -G^T \mathbf{y}$  to change  $\mathbf{e} = 0$  to  $\mathbf{e} \cdot \mathbf{y} = 0$ .
- ② But:  $\mathbf{e} \cdot \mathbf{y} = 0 \not\Rightarrow \mathbf{e} = 0$  even if  $\mathbf{y} \neq 0$ .
- ③ Redo: time-varying independent DG matrix  $Y' = -YG$  with  $Y\mathbf{e} = 0$ .
- ④  $Y\mathbf{e} = 0 \Rightarrow \mathbf{e} = 0$  if  $\det Y \neq 0$ .
- ⑤ DC  $\det Y \neq 0$  proves by dbx with Liouville:  $\det(Y)' = -\text{tr}(G)\det(Y)$
- ⑥ Continuous change of basis  $Y^{-1}$  balancing out motion of  $\mathbf{e}$ : constant!
- ⑦ Continuous change to new evolving variables is sound by DG. □

Derive  $[x' = f(x) \& Q](\mathbf{e})' = G\mathbf{e} \rightarrow (\mathbf{e} = 0 \rightarrow [x' = f(x) \& Q]\mathbf{e} = 0)$



Vectorial Darboux are DG

$$\frac{Q \rightarrow \mathbf{e}' = G\mathbf{e} \quad (G \in \mathbb{R}[x]^{n \times n})}{\mathbf{e} = 0 \rightarrow [x' = f(x) \& Q] \mathbf{e} = 0}$$



Proof Idea.

$$\begin{array}{c}
 * \\
 \text{R } (\mathbf{e})' = G\mathbf{e} \rightarrow -2\mathbf{e} \cdot (\mathbf{e})' \geq g(-\|\mathbf{e}\|^2) \\
 (\text{'})' \quad (\mathbf{e})' = G\mathbf{e} \rightarrow (-\|\mathbf{e}\|^2)' \geq g(-\|\mathbf{e}\|^2) \\
 \text{M } [x' = f(x) \& Q](\mathbf{e})' = G\mathbf{e} \rightarrow [x' = f(x) \& Q](-\|\mathbf{e}\|^2)' \geq g(-\|\mathbf{e}\|^2) \\
 \text{DBX } [x' = f(x) \& Q](\mathbf{e})' = G\mathbf{e}, -\|\mathbf{e}\|^2 \geq 0 \rightarrow [x' = f(x) \& Q] -\|\mathbf{e}\|^2 \geq 0 \\
 \text{M } [x' = f(x) \& Q](\mathbf{e})' = G\mathbf{e}, \mathbf{e} = 0 \rightarrow [x' = f(x) \& Q] \mathbf{e} = 0
 \end{array}$$

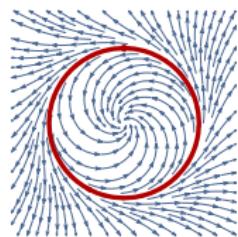
where  $g \stackrel{\text{def}}{=} 1 + \sum_{i=1}^n \sum_{j=1}^n G_{ij}^2$

1 + squared Frobenius □

Derive  $[x' = f(x) \& Q](\mathbf{e})' = G\mathbf{e} \rightarrow (\mathbf{e} = 0 \rightarrow [x' = f(x) \& Q] \mathbf{e} = 0)$

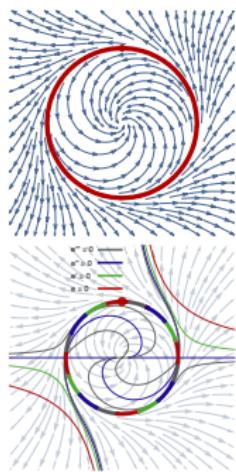
Vectorial Darboux are DG

$$\frac{Q \rightarrow \mathbf{e}' = G\mathbf{e}}{\mathbf{e} = 0 \rightarrow [x' = f(x) \& Q] \mathbf{e} = 0}$$



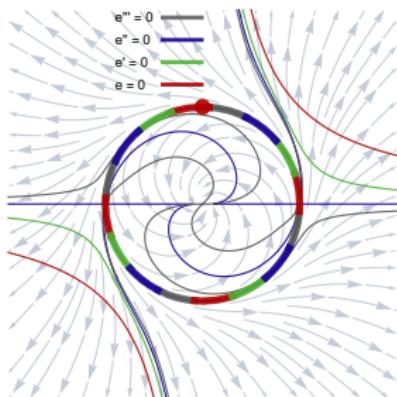
Vectorial Darboux are DG

$$\frac{Q \rightarrow \mathbf{e}' = G\mathbf{e}}{\mathbf{e} = 0 \rightarrow [x' = f(x) \& Q] \mathbf{e} = 0}$$



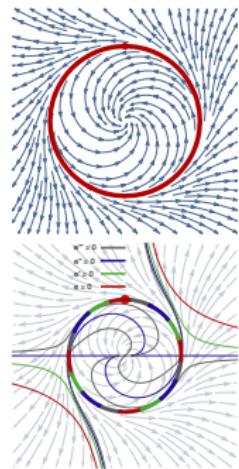
Differential Radical Invariants are DG

$$\frac{\Gamma, Q \rightarrow \bigwedge_{i=0}^{N-1} \mathbf{e}^{(i)} = 0 \quad Q \rightarrow \mathbf{e}^{(N)} = \sum_{i=0}^{N-1} g_i \mathbf{e}^{(i)}}{\Gamma \rightarrow [x' = f(x) \& Q] \mathbf{e} = 0}$$



Vectorial Darboux are DG

$$\frac{Q \rightarrow \mathbf{e}' = G\mathbf{e}}{\mathbf{e} = 0 \rightarrow [x' = f(x) \& Q]\mathbf{e} = 0}$$



Differential Radical Invariants are DG

$$\frac{\Gamma, Q \rightarrow \bigwedge_{i=0}^{N-1} \mathbf{e}^{(i)} = 0 \quad Q \rightarrow \mathbf{e}^{(N)} = \sum_{i=0}^{N-1} g_i \mathbf{e}^{(i)}}{\Gamma \rightarrow [x' = f(x) \& Q]\mathbf{e} = 0}$$

Proof Idea.

by vdbx with  $G = \begin{pmatrix} 0 & 1 & 0 & \dots & 0 \\ 0 & 0 & \ddots & \ddots & \vdots \\ \vdots & \vdots & \ddots & \ddots & 0 \\ 0 & 0 & \dots & 0 & 1 \\ g_0 & g_1 & \dots & g_{N-2} & g_{N-1} \end{pmatrix}$ ,  $\mathbf{e} = \begin{pmatrix} \mathbf{e} \\ \mathbf{e}^{(1)} \\ \mathbf{e}^{(2)} \\ \vdots \\ \mathbf{e}^{(N-1)} \end{pmatrix}$

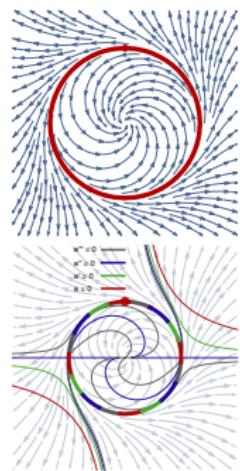
□

Vectorial Darboux are DG

$$\frac{Q \rightarrow \mathbf{e}' = G\mathbf{e}}{\mathbf{e} = 0 \rightarrow [x' = f(x) \& Q] \mathbf{e} = 0}$$

Differential Radical Invariants are DG

$$\frac{\Gamma, Q \rightarrow \bigwedge_{i=0}^{N-1} \mathbf{e}^{(i)} = 0 \quad Q \rightarrow \mathbf{e}^{(N)} = \sum_{i=0}^{N-1} g_i \mathbf{e}^{(i)}}{\Gamma \rightarrow [x' = f(x) \& Q] \mathbf{e} = 0} \quad N \text{ exists}$$



Vectorial Darboux are DG

$$\frac{Q \rightarrow \mathbf{e}' = G\mathbf{e}}{\mathbf{e} = 0 \rightarrow [x' = f(x) \& Q] \mathbf{e} = 0}$$

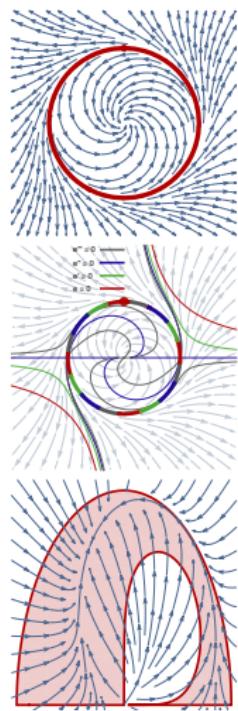
Differential Radical Invariants are DG

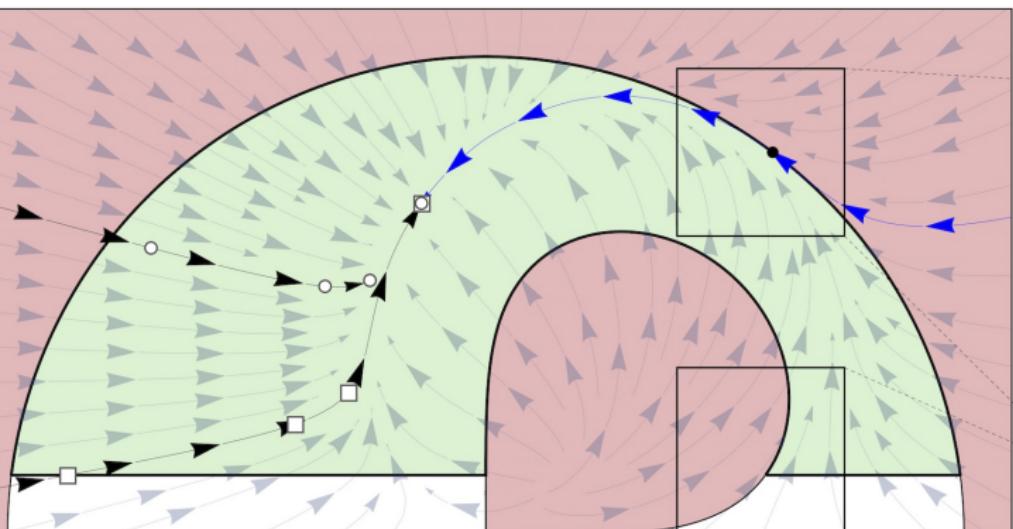
$$\frac{\Gamma, Q \rightarrow \bigwedge_{i=0}^{N-1} \mathbf{e}^{(i)} = 0 \quad Q \rightarrow \mathbf{e}^{(N)} = \sum_{i=0}^{N-1} g_i \mathbf{e}^{(i)}}{\Gamma \rightarrow [x' = f(x) \& Q] \mathbf{e} = 0} \quad N \text{ exists}$$

Semialgebraic Invariants are derived

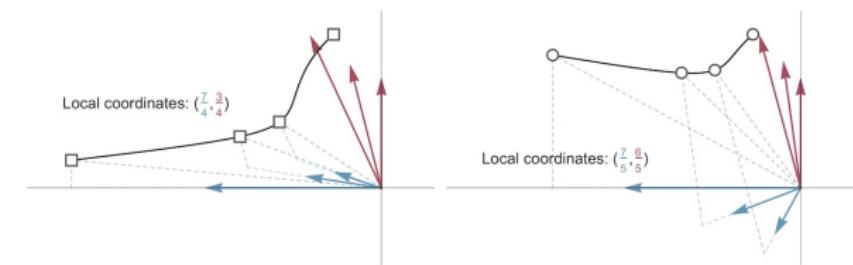
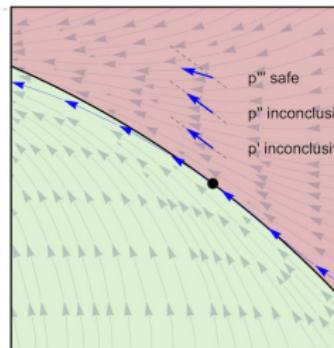
$$\frac{Q \rightarrow \mathbf{e}'^* \succcurlyeq 0 \quad Q \rightarrow \mathbf{e}^{(N)} = \sum_{i=0}^{N-1} g_i \mathbf{e}^{(i)}}{\mathbf{e} \succcurlyeq 0 \rightarrow [x' = f(x) \& Q] \mathbf{e} \succcurlyeq 0}$$

$$\begin{aligned} \mathbf{e}'^* \geq 0 &\equiv \mathbf{e} \geq 0 \wedge (\mathbf{e} = 0 \rightarrow (\mathbf{e}')'^* \geq 0) \\ \mathbf{e}'^* > 0 &\equiv \mathbf{e} > 0 \vee (\mathbf{e} = 0 \wedge (\mathbf{e}')'^* > 0) \end{aligned}$$

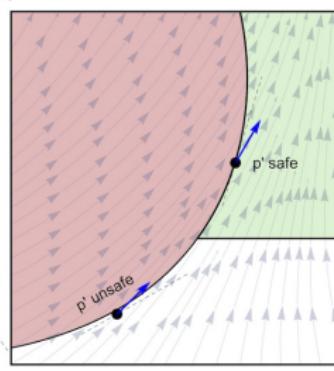




Proofs with higher Lie derivatives



Proofs use continuously changing basis ↪ to keep invariants at constant local coordinates



Sound and complete ODE invariance proofs

## Unique Solutions

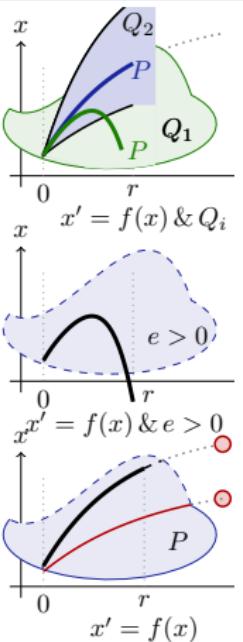
$$\langle x' = f(x) \& Q_1 \wedge Q_2 \rangle P \\ \leftrightarrow \langle x' = f(x) \& Q_1 \rangle P \wedge \langle x' = f(x) \& Q_2 \rangle P$$

## Continuous Existence

$$\langle x' = f(x) \& e > 0 \rangle x \neq y \leftrightarrow e > 0$$

## Real Induction

$$[x' = f(x)]P \leftrightarrow \forall y [x' = f(x) \& P \vee x = y] \\ (x = y \rightarrow P \wedge \langle x' = f(x) \& P \rangle x \neq y)$$



## Unique Solutions

$$\langle x' = f(x) \& Q_1 \wedge Q_2 \rangle P \\ \leftrightarrow \langle x' = f(x) \& Q_1 \rangle P \wedge \langle x' = f(x) \& Q_2 \rangle P$$

## Continuous Existence

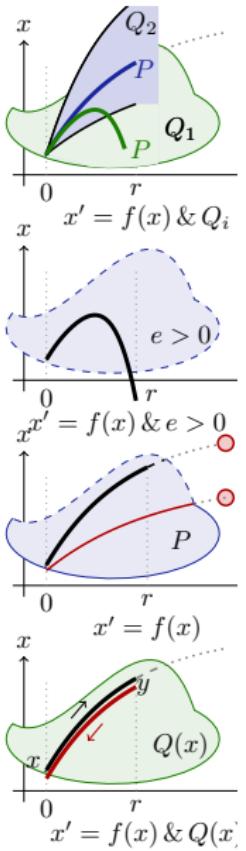
$$\langle x' = f(x) \& e > 0 \rangle x \neq y \leftrightarrow e > 0$$

## Real Induction

$$[x' = f(x)]P \leftrightarrow \forall y [x' = f(x) \& P \vee x = y] \\ (x = y \rightarrow P \wedge \langle x' = f(x) \& P \rangle x \neq y)$$

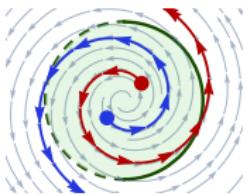
## Differential Adjoint

$$\langle x' = f(x) \& Q(x) \rangle x = y \leftrightarrow \langle y' = -f(y) \& Q(y) \rangle y = x$$



## Real Induction Rule

$$\frac{P \rightarrow \langle x' = f(x) \& P \rangle \bigcirc \quad \neg P \rightarrow \langle x' = -f(x) \& \neg P \rangle \bigcirc}{P \rightarrow [x' = f(x)]P}$$



$$\langle x' = f(x) \& P \rangle \bigcirc \equiv \langle y := x \rangle \langle x' = f(x) \& P \vee x = y \rangle_{x \neq y}$$

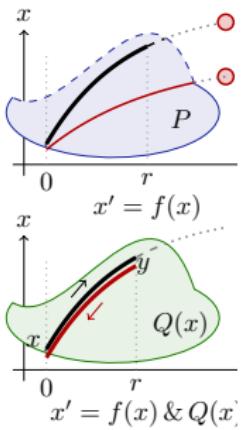
Local progress to  $P$

## Real Induction

$$[x' = f(x)]P \leftrightarrow \forall y [x' = f(x) \& P \vee x = y] \\ (x = y \rightarrow P \wedge \langle x' = f(x) \& P \rangle_{x \neq y})$$

## Differential Adjoint

$$\langle x' = f(x) \& Q(x) \rangle_{x=y} \leftrightarrow \langle y' = -f(y) \& Q(y) \rangle_{y=x}$$



## Local Progress Step

$$\begin{aligned} e > 0 \vee e = 0 \wedge \langle x' = f(x) \wedge e' \geq 0 \rangle \circ \\ \rightarrow \langle x' = f(x) \wedge e \geq 0 \rangle \circ \end{aligned}$$

Local Progress  $\geq$ 

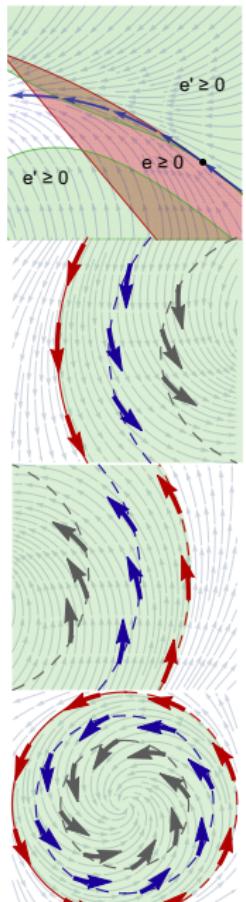
$$e'^* \geq 0 \rightarrow \langle x' = f(x) \wedge e \geq 0 \rangle \circ$$

Local Progress  $>$ 

$$e'^* > 0 \rightarrow \langle x' = f(x) \wedge e > 0 \rangle \circ$$

## Local Progress Semialgebraic

$$\langle x' = f(x) \wedge P \rangle \circ \leftrightarrow P'^*$$



## Local Progress Step

$$e > 0 \vee e = 0 \wedge \langle x' = f(x) \wedge e' \geq 0 \rangle \circ \\ \rightarrow \langle x' = f(x) \wedge e \geq 0 \rangle \circ$$

## Local Progress $\geq$

$$e'^* \geq 0 \rightarrow \langle x' = f(x) \wedge e \geq 0 \rangle \circ$$

## Local Progress $>$

$$e'^* > 0 \rightarrow \langle x' = f(x) \wedge e > 0 \rangle \circ$$

## Local Progress Semialgebraic

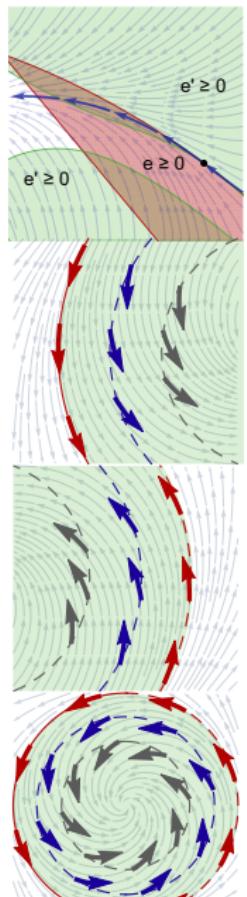
$$\langle x' = f(x) \wedge P \rangle \circ \leftrightarrow P'^*$$

$$e'^* \geq 0 \equiv e \geq 0 \wedge (e=0 \rightarrow (e')'^* \geq 0)$$

$$e'^* > 0 \equiv e > 0 \vee (e=0 \wedge (e')'^* > 0)$$

$$(P \wedge Q)'^* \equiv P'^* \wedge Q'^*$$

$$(P \vee Q)'^* \equiv P'^* \vee Q'^*$$



## Theorem (Algebraic Completeness)

(LICS'18,JACM'20)

dL calculus is a sound & complete axiomatization of algebraic invariants of polynomial differential equations. They are decidable with a derived axiom:

$$(\text{DRI}) \quad [x' = f(x) \& Q]e = 0 \leftrightarrow (Q \rightarrow e'^* = 0) \quad (Q \text{ open})$$

## Theorem (Semialgebraic Completeness)

(LICS'18,JACM'20)

dL calculus with RI is a sound & complete axiomatization of semialgebraic invariants of differential equations. They are decidable with derived axiom:

$$(\text{SAI}) \quad \forall x (P \rightarrow [x' = f(x)]P) \leftrightarrow \forall x (P \rightarrow P'^*) \wedge \forall x (\neg P \rightarrow (\neg P)^{**-})$$

Definable  $e'^*$  is short for all/significant Lie derivative w.r.t. ODE

Definable  $e^{**-}$  is w.r.t. backwards ODE  $x' = -f(x)$ . Also for  $P$ .

## Theorem (Analytic Completeness)

(LICS'18,JACM'20)

*dL calculus is a sound & complete axiomatization of analytic invariants of analytic differential equations.*

$$(\text{DRI}) \quad [x' = f(x) \& Q]e = 0 \leftrightarrow (Q \rightarrow e'^* = 0) \quad (Q \text{ open})$$

## Theorem (Semianalytic Completeness)

(LICS'18,JACM'20)

*dL calculus with RI is a sound & complete axiomatization of semianalytic invariants of differential equations.*

$$(\text{SAI}) \quad \forall x (P \rightarrow [x' = f(x)]P) \leftrightarrow \forall x (P \rightarrow P'^*) \wedge \forall x (\neg P \rightarrow (\neg P)^{**})$$

(S) Smooth function interpretations  $h : \mathbb{R}^k \rightarrow \mathbb{R}$

(P) Partial derivatives of  $h(y_1, \dots, y_k)$  have syntactic term representation  $\frac{\partial h}{\partial y_i}$

(R) Computable differential radicals: compute  $N, g_i$  for  $e^{(N)} = \sum_{i=0}^{N-1} g_i e^{(i)}$

## Definition (Noetherian Function)

$h : H \subseteq \mathbb{R}^k \rightarrow \mathbb{R}$  is *Noetherian function* iff  $h(y) = p(y, h_1(y), \dots, h_r(y))$  for a polynomial  $p$  and *Noetherian chain*  $h_1, \dots, h_r : H \subseteq \mathbb{R}^k \rightarrow \mathbb{R}$ , i.e., real analytic

$$\frac{\partial h_j}{\partial y_i}(y) = q_{ij}(y, h_1(y), \dots, h_r(y)) \text{ for some polynomial } q_{ij} \in \mathbb{R}[y, z]$$

Example:  $\frac{\partial \sin}{\partial y}(y) = \cos(y)$  and  $\frac{\partial \cos}{\partial y}(y) = -\sin(y)$  and  $\frac{\partial \exp}{\partial y}(y) = \exp(y)$

**Theorem** Noetherian functions satisfy SPR conditions.

⇒ Completeness for logic + differential equations with Noetherian functions.

(S) Smooth function interpretations  $h : \mathbb{R}^k \rightarrow \mathbb{R}$

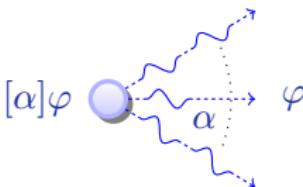
(P) Partial derivatives of  $h(y_1, \dots, y_k)$  have syntactic term representation  $\frac{\partial h}{\partial y_i}$

(R) Computable differential radicals: compute  $N, g_i$  for  $e^{(N)} = \sum_{i=0}^{N-1} g_i e^{(i)}$

- 1 Differential Dynamic Logic
  - Syntax
  - Axiomatization
  - Relative Completeness / ODE
- 2 Proofs for Differential Equations
  - Differential Invariants / Cuts / Ghosts
- 3 Completeness for Differential Equation Invariants
  - Darboux are Differential Ghosts
  - Derived Differential Radical Invariants
  - Real Induction
  - Derived Local Progress
  - Completeness for Invariants
  - Completeness for Noetherian Functions
- 4 Summary

differential dynamic logic

$$dL = DL + HP$$

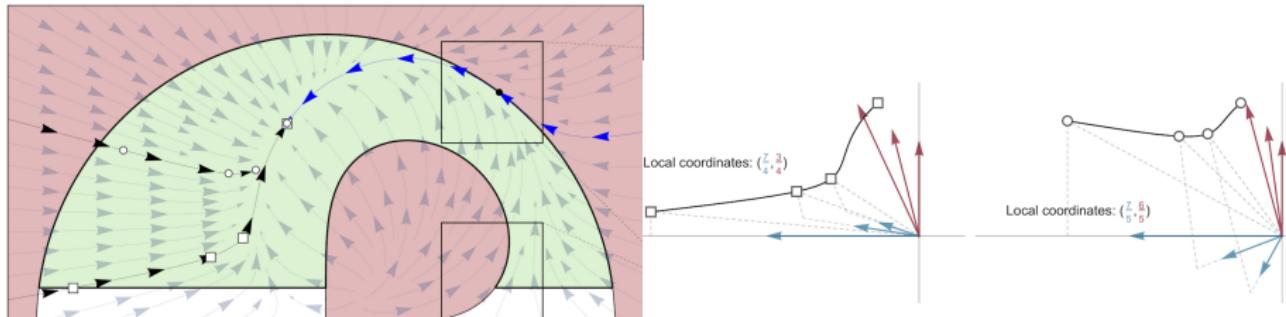
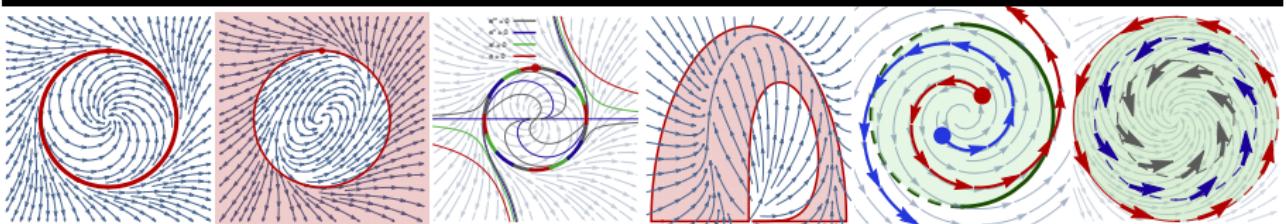
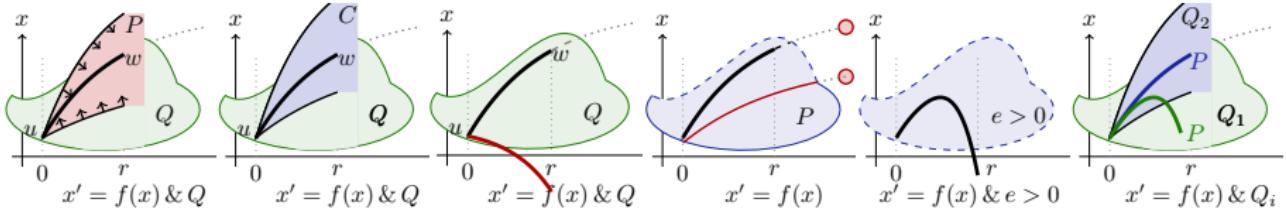


- 1 Poincaré: qualitative ODE
- 2 Complete axiomatization
- 3 Algebraic ODE invariants
- 4 Semialgebraic ODE invariants
- 5 Algebraic hybrid systems
- 6 Local ODE progress
- 7 Decide by dL proof/disproof
- 8 Uniform substitution axioms
- 9 Analytic extensions: Noetherian

## Properties

- |                          |                           |
|--------------------------|---------------------------|
| 1 MVT                    | 1 Differential invariants |
| 2 Prefix                 | 2 Differential cuts       |
| 3 Picard-Lind            | 3 Differential ghosts     |
| 4 $\mathbb{R}$ -complete | 4 Real induction          |
| 5 Existence              | 5 Continuous existence    |
| 6 Uniqueness             | 6 Unique solutions        |

Impressive power of differential ghosts



## I Part: Elementary Cyber-Physical Systems

1. Differential Equations & Domains
2. Choice & Control
3. Safety & Contracts
4. Dynamical Systems & Dynamic Axioms
5. Truth & Proof
6. Control Loops & Invariants
7. Events & Responses
8. Reactions & Delays

## II Part: Differential Equations Analysis

9. Differential Equations & Differential Invariants
10. Differential Equations & Proofs
11. Ghosts & Differential Ghosts
12. Differential Invariants & Proof Theory

## III Part: Adversarial Cyber-Physical Systems

- 13-16. Hybrid Systems & Hybrid Games

## IV Part: Comprehensive CPS Correctness



# Logical Foundations of Cyber-Physical Systems



André Platzer and Yong Kiam Tan.

Differential equation axiomatization: The impressive power of differential ghosts.

In Anuj Dawar and Erich Grädel, editors, *LICS*, pages 819–828, New York, 2018. ACM.

[doi:10.1145/3209108.3209147](https://doi.org/10.1145/3209108.3209147).



André Platzer and Yong Kiam Tan.

Differential equation invariance axiomatization.

*J. ACM*, 67(1):6:1–6:66, 2020.

[doi:10.1145/3380825](https://doi.org/10.1145/3380825).



André Platzer.

Differential dynamic logic for hybrid systems.

*J. Autom. Reas.*, 41(2):143–189, 2008.

[doi:10.1007/s10817-008-9103-8](https://doi.org/10.1007/s10817-008-9103-8).



André Platzer.

Logics of dynamical systems.

In LICS [10], pages 13–24.

[doi:10.1109/LICS.2012.13](https://doi.org/10.1109/LICS.2012.13).



André Platzer.

A complete uniform substitution calculus for differential dynamic logic.

*J. Autom. Reas.*, 59(2):219–265, 2017.

[doi:10.1007/s10817-016-9385-1](https://doi.org/10.1007/s10817-016-9385-1).



André Platzer.

The complete proof theory of hybrid systems.

In LICS [10], pages 541–550.

[doi:10.1109/LICS.2012.64](https://doi.org/10.1109/LICS.2012.64).



André Platzer.

Differential-algebraic dynamic logic for differential-algebraic programs.

*J. Log. Comput.*, 20(1):309–352, 2010.

[doi:10.1093/logcom/exn070](https://doi.org/10.1093/logcom/exn070).



André Platzer.

The structure of differential invariants and differential cut elimination.

*Log. Meth. Comput. Sci.*, 8(4:16):1–38, 2012.

[doi:10.2168/LMCS-8 \(4:16\) 2012](https://doi.org/10.2168/LMCS-8 (4:16) 2012).



André Platzer.

*Logical Foundations of Cyber-Physical Systems.*

Springer, Cham, 2018.

doi:10.1007/978-3-319-63588-0.



*Logic in Computer Science (LICS), 2012 27th Annual IEEE Symposium on*, Los Alamitos, 2012. IEEE.