

# D-finite functions (univariate case)

Defn: A function  $f(x)$  is called D-finite (differentially finite) if it satisfies a LODE with polynomial coeffs: or: holonomic

$$p_r(x)f^{(r)}(x) + \dots + p_1(x)f'(x) + p_0(x)f(x) = 0$$

with  $p_0, \dots, p_r \in K[x]$  (not all zero)

Examples:  $x^n$ , const.,  $\sin$ ,  $\exp$ ,  $\dots$

not D-finite:  $\tan$ ,  $\Gamma(x)$



Features:

- finitely many initial consts.  $\rightarrow$  fin. amount of data
- rich class (show slide)
- closure properties

Thm: If  $f(x)$  and  $g(x)$  are D-finite, then also  $f(x) + g(x) = h(x)$  is DF.

Proof: Let  $r, s$  denote the orders of  $f, g$ , resp. To show:

$$\exists u \in \mathbb{N}, p_0, \dots, p_u \in K[x]: p_u(x)h^{(u)}(x) + \dots + p_0(x)h(x) = 0.$$

$$\Leftrightarrow p_u(x)(f^{(u)}(x) + g^{(u)}(x)) + \dots + p_0(x)(f(x) + g(x)) = 0$$

Using the LODE for  $f$ :  $f^{(u)}(x) = \{ \cdot \} \cdot f^{(u-1)}(x) + \dots + \{ \cdot \} \cdot f'(x) + \{ \cdot \} \cdot f(x)$   
(and analogously for  $g$ ) — works for any derivative

$$\Rightarrow l_0(p_0, \dots, p_u) \cdot f(x) + l_1(p_0, \dots, p_u) \cdot f'(x) + \dots + l_{r+s}(p_0, \dots, p_u) \cdot f^{(r+s)}(x)$$

$$+ l_r(p_0, \dots, p_u) \cdot g(x) + \dots + l_{r+s-1}(p_0, \dots, p_u) \cdot g^{(s-1)}(x) = 0$$

where each  $l_j$  is a linear poly in the  $p_i$ :  $l_j = \sum_{i=0}^u c_{i,j}(x) \cdot p_i(x)$ .

The linear system  $l_0 = \dots = l_{r+s-1} = 0$  is guaranteed to have a nontrivial solution if  $u = r+s$  is chosen.  $\square$

Hint: there may exist a solution for smaller  $u$ , hence try and loop.

Thm: If  $f(x)$  and  $g(x)$  are D-finite, then also

$f(x) \cdot g(x)$  (proof: analogous, exercise!)

$f'(x)$  (or more general: any diff. poly in  $f, f', \dots$ )

$f(\alpha(x))$  where  $\alpha(x)$  is an algebraic function;  $p(x, \alpha) = 0$  for  $p \in K[x, y]$

$\int f(x) dx$  (proof: replace  $f^{(t)}$  by  $f^{(t+1)}$ )

Platz für  
Ihre Ideen.

- Operator notation: Let  $D_x$  denote the differentiation w.r.t.  $x$ , i.e.,  
 $D_x(f(x)) = f'(x)$ ,  $D_x^2(f(x)) = f''(x)$ , etc.  $D_x^0(f(x)) = f(x)$
- Let  $K(x)\langle D_x \rangle$  denote the polynomial ring in  $D_x$  with coeffs in  $K(x)$ .  
 It is not commutative  $D_x \cdot x = x \cdot D_x + 1$  (Leibniz rule)  
 more general:  $D_x \cdot r(x) = r(x) \cdot D_x + r'(x)$
  - $f(x)$  is  $D$ -finite  $\Leftrightarrow \exists L \in K(x)\langle D_x \rangle \setminus \{0\} : L(f(x)) = 0$

Fact: Let  $L_1, L_2 \in K(x)\langle D_x \rangle \setminus \{0\}$  annihilate  $f, g$  resp.,  $L_1(f(x)) = 0$ . Then  

- $L \cdot D_x$  annihilates  $\int f(x) dx$
- lcm ( $L_1, L_2$ ) annihilates  $f+g$  (actually  $c_1 f + c_2 g$  for arb. const.  $c_1, c_2$ )  
 $\Rightarrow L$  Proof:  $L = M_1 L_1 = M_2 L_2$  for certain  $M_1, M_2 \in K(x)\langle D_x \rangle$
- If  $f$  satisfies  $L(f) = g$  for some  $D$ -finite  $g$ , then  $f$  is  $D$ -finite.  
 Proof: Assume  $M(g) = 0$ . Then  $(M \cdot L)(f) = 0$ .

AEZDE  $\rightarrow$  Corollary of previous Thm (with  $f(x) = x$ )

Thm: Let  $f(x)$  be an algebraic function. Then  $f$  is  $D$ -finite.

Proof: Let ~~possibly~~  $m \in K[x, y]$  be the minimal polynomial of  $f$ ,  
 i.e.,  $m(x, f(x)) = 0$  and  $m$  irreducible.

$$\hookrightarrow a_d(x)(f(x))^d + a_{d-1}(x)(f(x))^{d-1} + \dots + a_1(x)f(x) + a_0(x) = 0$$

$$\text{Differentiate: } a'_d f^d + d \cdot a_d f^{d-1} f' + \dots + a'_1 f + a'_0 + a'_0 = 0$$

$$\Rightarrow f' = \frac{-(a'_d f^d + \dots + a'_1 + a'_0)}{d \cdot a_d f^{d-1} + \dots + a'_1} =: \frac{q(x, f)}{r(x, f)}$$

Note:  $\gcd(m, r) = 1$ , hence by EEA  $\exists u, v \in K(x)[f] : u \cdot m + v \cdot r = 1$

We get  $v(x, f) \cdot r(x, f) = 1 \pmod{m(x, f)}$ , hence

$$f' = \frac{q(x, f) \cdot v(x, f)}{r(x, f) \cdot v(x, f)} = q(x, f) \cdot v(x, f) = c_{1,d-1} f^{d-1} + \dots + c_{1,1} f + c_{1,0} \pmod{m}$$

Differentiate again:

$$f'' = c_{2,d-1} f^{d-1} + \dots + c_{2,1} f + c_{2,0}, \text{ etc.}$$

It follows that  $f$  satisfies a LODE with polynomial coeffs.  
 of order at most  $d$ .

Def.: A sequence  $(a_n)_{n \in \mathbb{N}}$  is called P-recursive | or: holonomic if it satisfies a linear recurrence equation | or: D-finite with polynomial coefficients:

$$P_r(n) \cdot a_{n+r} + \dots + P_1(n) a_{n+1} + P_0(n) a_n = 0$$

with  $P_0, \dots, P_r \in K[n]$  (not all zero).

Examples: Fibonacci, polynomials, hypergeometric,  $H_n$ , ortho. polys  
not P-rec: prime numbers

Features:

- finitely many initial values  $\rightarrow$  finite amount of data
- rich class (show slide)
- closure properties

Thm: If  $a_n$  and  $b_n$  are P-recursive, then also the following are P-rec.

- $a_n + b_n$
- $a_n \cdot b_n$
- $a_{cn+d}$  for integers  $c, d$
- $\sum_n a_n$  (indefinite sum, i.e.,  $c_n$  s.t.  $c_{n+1} - c_n = a_n$ )

Operator notation:

- $S_n$  denotes the forward shift operator, i.e.,  $S_n(a_n) = a_{n+1}, S_n^2(a_n) = a_{n+2}$
- $K(n)\langle S_n \rangle$ : ring of all linear recurrence operators with coeffs in  $K(n)$   
commutation rule:  $S_n \cdot n = (n+1) \cdot S_n$ , more gen.:  $S_n \cdot r(n) = r(n+1) \cdot S_n$

Facts: see D-finite

Thm: A function  $f(x)$  is D-finite if and only if the sequence of its Taylor coefficients is P-recursive. Stated differently:

A sequence is P-rec. iff its generating function is Definite.

Proof: Let  $f(x) = \sum_{n=0}^{\infty} a_n x^n$ . Then  $f'(x) = \sum_{n=1}^{\infty} n \cdot a_n x^{n-1}$ ,  $f^{(i)}(x) = \sum_{n=i}^{\infty} (n-i+1)_i a_n x^{n-i}$

$$\text{LODE: } \sum_{i=0}^r \sum_{j=0}^d p_{i,j} x^j f^{(i)}(x) = 0$$

$$= \sum_{n=0}^{\infty} (n+1)_i a_{n+i} x^n$$

$$\Rightarrow \sum_{i=0}^r \sum_{j=0}^d \sum_{n=0}^{\infty} p_{i,j} (n+1)_i a_{n+i} x^{n+j} = 0$$

$$\Rightarrow \sum_{i=0}^r \sum_{j=0}^d \sum_{n=j}^{\infty} p_{i,j} (n-j+1)_i a_{n-j+i} x^n \Rightarrow \sum_{i=0}^r \sum_{j=0}^d p_{i,j} (n-j+1)_i a_{n+j-i} = 0$$

↑ for  $i \leq d$