

Manuel Kauers · Institute for Algebra · JKU

Monday	Tuesday	Wednesday	Thursday	Friday
Intro- duction	Binomial Summation	Creative Telescoping	D-finite univariate	Chyzak's algorithm
Rational Integration Theory	Sister Celine Theory	Gosper's algorithm	D-finite multivariate	Example Session
Rational Integration Coding	Sister Celine Coding	Zeilberger's algorithm	Advanced Closure Properties	Conclusion

# Some points to remember:

- What is a telescoper?
- What is it good for?
- How can it be computed?

# What did we not cover in this course?

- Liouvillean functions and  $\Pi\Sigma$  expressions
- Reduction-based creative telescoping for D-finite functions

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The field K together with such a D is called a differential field.

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**Example:** On  $K = \mathbb{Q}(t_1, t_2, t_3, t_4)$  we can define a derivation via

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Example: 
$$\int \frac{1}{1 + \exp(x)} = x - \log(1 + \exp(x))$$



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We can also use this for evaluating definite integrals of liouvillean functions.
	hypergeometric	liouvillean	
	summation	integration	
indefinite	Gosper	Risch	
definite (CT)	Zeilberger	Raab	

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# Example: $\sum_{k=1}^{n} \frac{2^{k} - \sum_{i=1}^{k} \frac{1}{k}}{k! + \sum_{i=1}^{k} \frac{1}{k^{2}}}$

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A difference field is a field K together with an automorphism  $\sigma \colon K \to K.$ 

**Example:** On  $K = \mathbb{Q}(t_1, t_2, \dots)$  we can define  $\sigma$  via

$$\begin{split} &\sigma(t_1) = t_1 + 1 & t_1 \sim n \\ &\sigma(t_2) = 2t_2 & t_2 \sim 2^n \\ &\sigma(t_3) = (t_1 + 1)t_3 & t_3 \sim n! \\ &\sigma(t_4) = t_4 + \frac{1}{t_1 + 1} & t_4 \sim \sum_{k=1}^n \frac{1}{k}, \quad \text{etc.} \end{split}$$

Def. A difference field  $K=C(t_1,\ldots,t_d)$  is called  $\Pi\Sigma$  if the difference subfield  $C(t_1,\ldots,t_{d-1})$  is  $\Pi\Sigma$  and

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Karr's algorithm solves the summation problem in such fields:

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Example: 
$$\sum_{k=1}^{n} \sum_{i=1}^{k} \frac{1}{k} = (n+1) \sum_{k=1}^{n} \frac{1}{k} - n$$

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Schneider uses it to do creative telescoping and lots of other things.

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	summation	integration	summation
indefinite	Gosper	Risch	Karr
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#### What did we not cover in this course?

- Liouvillean functions and  $\Pi\Sigma$  expressions
- Reduction-based creative telescoping for D-finite functions

#### Recall:

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- Zeilberger-like algorithms are based on an indefinite summation/integration algorithm
- Apagodu-Zeilberger-like algorithms are based on an ansatz for telescoper and certificate and solving a linear system
- Reduction-based algorithms are based on extracting maximal summable/integrable parts

Reduction-based creative telescoping for D-finite functions Example: Hermite reduction breaks a given  $f \in C(x, y)$  into

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- Hyperexponential functions (Bostan, Chen, Chyzak, Li, Xin)
- Hypergeometric terms (Abramov, Petkovšek; Chen, Huang, Li, Kauers)
- D-finite functions

(Bostan, Brochet, Chen, Du, van Hoeij, van der Hoeven, Lairez, Kauers, Koutschan, Salvy, Wang)

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$$\begin{split} &\sum_{k=1}^{n-1} \frac{k^4 + 5k^3 - k^2 - 5k - 2}{k^2(k+1)^3(k+2)} \binom{2k}{k} \\ &= -\frac{n^2 - 6}{6n^2(n+1)} \binom{2n}{n} + \frac{5}{6} + \sum_{k=1}^{n-1} \frac{1}{2(k+1)} \binom{2k}{k}. \end{split}$$

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We obtain reduction-based creative telescoping algorithms.

These techniques are still subject of ongoing research.

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# What remains to be done in the future?

J Syst Sci Complex (2017) 30: 154-172

#### Some Open Problems Related to Creative Telescoping\*

CHEN Shaoshi · KAUERS Manuel

DOI: 10.1007/s11424-017-6202-9 Received: 14 September 2016 / Revised: 30 October 2016 ©The Editorial Office of JSSC & Springer-Verlag Berlin Heidelberg 2017

Abstract Creative telescoping is the method of choice for obtaining information about definite sums or integrals. It has been intensively studied since the early 1990s, and can now be considered as a classical technique in computer algebra. At the same time, it is still a subject of ongoing research. This paper presents a selection of open problems in this context. The authors would be curious to hear about any substantial progress on any of these problems.

Keywords Computer algebra, creative telescoping, differential algebra, linear operators, ore algebras, symbolic integration, symbolic summation.

#### 1 Introduction

Summation and integration problems arise in all areas of mathematics, especially in discreto mathematics, special functions, combinatorics, engineering, and physics. Nowadays, many o these problems are solved using computer algebra. The number of applications of summation • Reduction-based telescoping for further function classes.

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- The inverse problem of definite summation/integration.

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- Multivariate indefinite summation/integration.
- Integration of D-algebraic functions.
- The inverse problem of definite summation/integration.
- Software that can handle problems out of reach of available code.