

Hypergeometric Summation: Gosper's Algorithm

A sequence $H(n)$ is hypergeometric over $K(n)$ if

$$\frac{H(n+1)}{H(n)} \in K(n).$$

Examples: $\frac{1}{n}$, n^2+1 , 2^n , $n!$, $\Gamma(\alpha+n)$, ...

Two hypergeometric terms $H_1(n)$, $H_2(n)$ are said to be similar if $H_1/H_2 \in K(n)$, denoted by $H_1 \sim H_2$.

problem Given a hypergeometric $H(n)$, decide

whether there exists another hypergeometric $G(n)$

s.t.

$$H(n) = G(n+1) - G(n) \stackrel{\Delta}{=} \Delta_n(G(n)).$$

If such a $G(n)$ exists, we say that $H(n)$ is hypergeometric

summable

Remark

If $H(n) = \Delta_n(G(n))$ for some hypergeom.

$G(n)$, then $G(n+1) = g(n)G(n)$ for some rational

$g \in K(n)$. Thus $H(n) = (g(n) - 1)G(n)$

which means $H(n)$ and $G(n)$ are similar.

To decide the existence of $G(n)$, we make an ansatz

$$G(n) = y(n) \cdot H(n) \quad \text{Assume } \frac{H(n+1)}{H(n)} = f(n) \in K(n)$$

$$\begin{aligned} H(n) = \Delta_n(G(n)) &\Leftrightarrow H(n) = y(n+1)H(n+1) - y(n)H(n) \\ &= (y(n+1)f(n) - y(n))H(n) \end{aligned}$$

\Rightarrow (*) $| = f(n)y(n+1) - y(n)$ has a rational solution in $K(n)$

Gosper (1978) gave a nice way to guess the denominator of $y(n)$ and a degree bound for the numerator of $y(n)$.

DEF (Gosper form) Let $f \in K(n)$. We call the triple $(P, Q, R) \in K(n)^3$ a Gosper form of f if

$$f = \frac{P(n+1)}{P(n)} \frac{Q(n)}{R(n)}$$

where $\gcd(Q(n), r(n+i)) = 1 \quad \forall i \in \mathbb{N}$

Theorem (*) Any rational function has a Gosper form.

proof Let $f = \frac{a}{b}$ with $\gcd(a, b) = 1$. If

$g = \gcd(a(n), b(n+j_0)) \neq 1$ for some $j_0 \in \mathbb{N}$
then $g(n) \mid a(n)$ and $g(n) \mid b(n+j_0)$

$$\textcircled{2} \Rightarrow g(n-j_0) \mid b(n)$$

Let $a = g(n) \bar{a}$ $b = g(n-j_0) \bar{b}$

$$\frac{a(n)}{b(n)} = \frac{g(n) \bar{a}}{g(n-j_0) \bar{b}} = \frac{g(n)}{g(n-1)} \frac{g(n-1)}{g(n-2)} \dots \frac{g(n-j_0+1)}{g(n-j_0)} \frac{\bar{a}}{\bar{b}}$$

$$= \frac{P_1(n+1)}{P_1(n)} \frac{\bar{a}}{\bar{b}} \quad \text{with } P_1 = g(n-1) \dots g(n-j_0)$$

We can iterate the process for \bar{a}/\bar{b} until it satisfies the GCD condition (*).

Algorithm (Gosper 1978)

• Equation to solve:
 $1 = f(n) y(n+1) - y(n)$

Input: A hypergeometric $H(n)$

output: No or $y(n)$ s.t. $H(n) = \Delta_n (y(n) H(n)) \in K(n)$

Step 1: Compute a Gosper form of f

$$f = \frac{p(n+1)}{p(n)} \frac{q(n)}{r(n)}$$

Step 2 Make an ansatz: $y(n) = \frac{r(n-1)}{p(n)} z(n)$

$$\text{Then } 1 = f(n) y(n+1) - y(n) \Leftrightarrow 1 = \frac{p(n+1)}{p(n)} \frac{q(n)}{r(n)} \frac{r(n)}{p(n+1)} z(n+1) - \frac{r(n-1)}{p(n)} z(n)$$

$$\Leftrightarrow p(n) = q(n) z(n+1) - r(n-1) z(n)$$

(*) (*) Gosper's equation

Gosper's Key Lemma

The equation $1 = f(n)y(n+1) - y(n)$

has a rational solution

\Leftrightarrow the Gosper equation

(**) $P(n) = z(n)z(n+1) - r(n-1)z(n)$
has ^{only} a polynomial solution

Proof " \Leftarrow " clear

" \Rightarrow " If $1 = f(n)y(n+1) - y(n)$ has a rational solution then (**) also has a rational solution, say

$$z(n) = \frac{a(n)}{b(n)} \in K(n) \text{ with } \gcd(a, b) = 1$$

$$\Rightarrow P(n) = z(n) \frac{a(n+1)}{b(n+1)} - r(n-1) \frac{a(n)}{b(n)}$$

Claim $b(n) \in K^*$. otherwise, assume that $b \notin K^*$

Then
$$P(n) \underbrace{b(n)}^m \underbrace{b(n+1)}^{\tilde{m}} = z(n) a(n+1) \tilde{b}(n) - r(n-1) a(n) \tilde{b}(n+1)$$

Take any irreducible polynomial $u(n)$ s.t. $u(n) \mid b(n)$

Then \exists maximal $j \in \mathbb{N}$ s.t.

$$u(n+j) \mid b(n) \quad \text{but} \quad u(n+j+1) \nmid b(n)$$

$$\gcd(a(n), b(n)) = 1, \Rightarrow \begin{array}{l} u(n) \mid b(n) \Rightarrow u(n) \nmid a(n) \\ u(n+j) \mid b(n) \Rightarrow u(n+j) \nmid a(n) \end{array}$$

$$u(n+j) \mid b(n) \Rightarrow u(n+j+1) \nmid a(n+1)$$

$$\Rightarrow u(n+j+1) \mid b(n+1) \Rightarrow u(n+j+1) \mid z(n) a(n+1) b(n)$$

$$\Rightarrow \boxed{u(n+j+1) \mid z(n)}$$

On the other hand, \exists maximal $i \in \mathbb{N}$ s.t.

$$u(n-i) \mid b(n) \text{ but } u(n-i-1) \nmid b(n)$$

$$\Rightarrow u(n-i) \nmid b(n+1)$$

$$u(n-i) \mid b(n) \Rightarrow u(n-i) \mid r(n-1) a(n) b(n+1)$$

$$\Rightarrow \boxed{u(n-i) \mid r(n-1)} \quad \left(\begin{array}{l} u(n-i) \nmid a(n) \\ \text{since } \gcd(a(n), b(n)) = 1 \end{array} \right)$$

So we now get

$$\left\{ \begin{array}{l} u(n+i+1) \mid z(n) \\ u(n-i+1) \mid r(n) \end{array} \right. \Rightarrow \gcd(z(n), r(n+i+j)) \neq 1$$

$i, j \in \mathbb{N}$

→ ← contradiction!!!

Step 3 Find a polynomial solution of

$$p(n) = z(n) z(n+1) - r(n-1) z(n)$$

$$\Leftrightarrow p(n) = z(n) (z(n+1) - z(n)) + (z(n) - r(n+1)) z(n)$$

$$= z(n) \Delta(z(n)) + (z(n) - r(n+1)) z(n)$$

Now we can estimate the degree bound of $z(n)$

and then apply the method of undetermined coefficients

we finally leads to solving a linear system over K .

Estimating the degree bound for the equation

$$a(n) = b(n)\Delta(z(n)) + c(n)z(n) \quad (*)$$

$$z = z_d x^d + z_{d-1} x^{d-1} + \dots + z_0$$

$$\Delta(z) = d z_d x^{d-1} + \text{lower terms}$$

Case 1 $\deg_x b \leq \deg_x c$

Then $d = \deg_x a - \deg_x c$

Case 2 $\deg_x b > \deg_x c + 1$

Then $d = \deg_x a - \deg_x b + 1$

Case 3 $\deg_x b = \deg_x c + 1 = p$

$$b = b_p x^p + \dots + b_0 \quad c = c_{p-1} x^{p-1} + \dots + c_0$$

$$b\Delta(z) + cz = (b_p d z_d + c_{p-1} z_d) x^{p+d-1} + \text{lower terms}$$

Then either $d = -\frac{c_{p-1}}{b_p}$ or $p+d-1 = \deg_x a$

If $-\frac{c_{p-1}}{b_p} \notin \mathbb{N}$, then $d = \deg_x a - p - 1$

Otherwise $d = \max \left\{ -\frac{c_{p-1}}{b_p}, \deg_x a - p - 1 \right\}$

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Example $H = \frac{\binom{m}{k}}{\binom{n}{k}}$

Claim $H = \Delta_k(G)$

$$G = \frac{n-k+1}{m-n+1} \cdot H(k)$$

We now compute G by Gosper's Algorithm.

Step 1 $\frac{H(k+1)}{H(k)} = \frac{m-k}{n-k} = \frac{p(k+1)}{p(k)} \frac{q(k)}{r(k)}$

$$\Rightarrow p(k) = 1 \quad q = m-k \quad r = n-k$$

Step 2 Finding a polynomial solution of the equation

$$1 = (m-k)z(k+1) - (n-k+1)z(k)$$

$$= (m-k)\Delta_k(z(k)) + (m-n-1)z(k)$$

$$\Rightarrow \deg_k(z) = 0 \Rightarrow z(k) = \frac{1}{m-n-1}$$

$$\Rightarrow y(k) = \frac{r(k-1)}{p(k)} z(k) = \frac{n-k+1}{m-n-1}$$

Then $\sum_{k=0}^m H(k) = \frac{n+1}{n-m+1}$

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